

A NOTE ON THE IMPROVEMENT OF THE MAXIMUM INDEPENDENT SET'S APPROXIMATION RATIO

Vangelis TH. PASCHOS

*Lamsade, Universite Paris-Dauphine
Place du Maréchal De Lattre de Tassigny, 75775 Paris Cedex 16, France*

Abstract: We present an $O(n^{3.5})$ approximation algorithm for maximum independent set problem achieving, in unweighted case, a worst case approximation ratio strictly greater than $(2/\Delta)$, for a graph of order n and maximum vertex degree Δ . The best known ratio was, up to now, equal to $2/\Delta$ and its improvement is considered as an interesting open problem.

Keywords: Approximation algorithms, analysis of algorithms, combinatorial problems, independent set, computational complexity.

The aim of this note is to propose a new approximation algorithm for maximum independent set problem which, for the unweighted case, guarantees an approximation ratio better (greater) than $2/\Delta$ which, up to now, is the best known ratio [3, 5].

Let $G = (V, E)$ be a graph of order n ; an independent set V' is a subset of V such that whenever $\{v_i, v_j\} \subseteq V'$, $v_i v_j \notin E$, and the maximum independent set problem (IS) is to find an independent set of maximum size; in what follows, we denote by $\alpha(G)$ the stability number (cardinality of a maximum independent set) of G .

Nemhauser and Trotter prove, in [4], that the LP-relaxation¹ IS_r of IS has the semi-integral property, i.e., each basic feasible optimal solution F of IS_r assigns to the variables values drawn from the set $\{0, 1/2, 1\}$ (let us denote by $v(F)$ the value of F , and consequently, the optimal value of IS_r); next, starting from this fact, they prove that if S is the set of vertices corresponding to variables valued by 1 in F , then there exists a maximum independent set in G that contains S . This elegant result, is at the

¹ We replace the integrality constraints on the IS-solution by positivity ones.

origin of the very interesting Hochbaum's IS-algorithm [3].²

Here, we show how we can exploit further the result of [4] in such a way that, in combination with the natural greedy heuristic for IS, we obtain a better approximation bound for IS, i.e., an approximation ratio strictly greater than $2/\Delta$. The algorithm we propose uses as procedure an algorithm, proposed by Bourjolly et al. in [1], that decides if, in a graph G , $\nu(F) = \alpha(G)$, and if this is true, it finds a maximum independent set in G in $O(n^3)$; in fact, it is proved there, that, in a graph G , the property "stability number equal size of a minimum edge covering", property which defines the class of König-Egervary graphs (KE-graphs), is equivalent to the property " $\nu(F) = \alpha(G)$ ".

begin

$S' \leftarrow \emptyset; V_C \leftarrow V$

while the variables of F for $G[V_C]$ are not all equal to $1/2$ **do**

 solve IS_r in order to define sets $S, \bar{S}; V_C \leftarrow V_C \setminus (S \cup \bar{S}); S' \leftarrow S' \cup S$

od

 call the algorithm of [1] to decide if $G[V_C]$ is KE;

if $G[V_C]$ is KE **then**

 use the algorithm of [1] to solve (optimally) IS on $G[V_C]$;

 let S_C be the IS-solution for $G[V_C]$; $S' \leftarrow S' \cup S_C$

exit

fi

repeat

$v_j \leftarrow \operatorname{argmin}_{v_i \in V_C} \{\delta_i\}; S' \leftarrow S' \cup \{v_j\}; V_C \leftarrow V_C \setminus (\{v_j\} \cup \Gamma(v_j));$

 delete all edges incident to $\{v_j\} \cup \Gamma(v_j)$;

 update the degrees of the vertices in V_C

until $V_C = \emptyset$

end.

Algorithm 1. A modified greedy IS-algorithm; the solution constricted is denoted by S' .

In what follows, we consider a connected graph³ $G = (V, E)$ of order n ; for a $V' \subseteq V$, we denote by $G[V']$ the subgraph of G induced by V' ; for every vertex v_i of G , we denote by $\Gamma(v_i)$ the set of its neighbours, by δ_i the degree of v_i , and by Δ the quantity $\max_{v_i \in V} \delta_i$; also, let us denote by S and \bar{S} the subsets of V to the corresponding variables of which F assigns values 1 and 0, respectively.

² We have to mention that Hochbaum is the first author who has shown how we can use Nemhauser-Trotter's method to obtain efficient bounds for combinatorial problems.

³ If the graph is not connected, then the result of the theorem remains true by considering every connected component of the graph.

THEOREM. Algorithm 1 is an $O(n^{3.5})$ approximation IS-algorithm achieving a worst case approximation ratio greater than or equal to $(2/\Delta)[1 + 1/(n-2)]$.

Let us suppose that the **while** loop of Algorithm 1 is repeated t times. Let us denote by $\alpha(G_j)$, $j = 1, \dots, t$, the stability number of the surviving graph at the beginning of the i th iteration ($\alpha(G_j) = \alpha(G)$); of course, this graph is a graph $G[V_C^j]$ where, for purposes of facility, we denote by V_C^j the set V_C produced at the end of the j -1th iteration. Also, for an iteration j , let us denote by S_j the set - and by σ_j the cardinality - of vertices, the corresponding variables of which are assigned by value 1 in F ; it is easy to see, from the constraint-set of the integer-linear program of IS, that the set $\cup_{i \leq j} S_i$ is an independent set of G , so does the union of this set with an independent set of the graph surviving after the t executions of the **while** loop.

Then, based on the result of [4], for $j = 1, \dots, t$, there exists a maximum independent set of $G[V_C^j]$ containing S_j ; we so obtain a t -line system of equations

$$\alpha(G[V_C^j]) = \alpha(G[V_C^{j+1}]) + \sigma_j, \quad j = 1, \dots, t$$

where $G[V_C^1] = G$, and $G[V_C^{t+1}]$ is the final graph, all the vertices of which have their corresponding variables valued by 1/2 in F ; let us denote $G[V_C^{t+1}]$ by G_{t+1} .

By summing the lefthand and righthand sides of the equations of the above system, we obtain

$$\alpha(G) = \sum_{j=1}^t \sigma_j + \alpha(G_{t+1}) \quad (1)$$

If, after the execution of the **while** loop of Algorithm 1, the **if** condition is executed, then the optimal solution (and, hence, an approximation ratio equal to 1) is found for G .

Plainly, since IS in KE-graphs is polynomial, the value of the set S_C equals $\alpha(G_{t+1})$ (where S_C is the set appearing in the third line of the **if** condition); so, the solution S' produced by Algorithm 1 just before the **exit** instruction has value

$\sum_{j=1}^t \sigma_j + \alpha(G_{t+1})$ and, by expression (1), this cardinality is equal to the optimal one.

Let us now estimate the approximation performance of Algorithm 1 (**repeat** loop) on G_{t+1} (the execution of this loop survenes in the case where G_{t+1} is not KE).

Let us denote by n_{t+1} the order of G_{t+1} ; we also denote by S'_{t+1} the subset of vertices added in S' during the execution of the second line of the **repeat** loop.

So, during the first iteration of the **repeat** loop, since the minimum-degree vertex is chosen, at most $\Delta + 1$ vertices of V_C^{t+1} are deleted, one among them being included in S'_{t+1} (or, equivalently, in S'). Next, during ulterior iterations of the **repeat**

loop (and because of the facts that (i) the minimum degree vertex is always chosen and (ii) the graph G_{t+1} is connected³), the degree of the chosen vertices will be at most $\Delta - 1$; so, at most Δ vertices will be deleted from V_C^{t+1} ; this is due to the fact that, by the connexity of the graph G_{t+1} , there always exists at least one vertex, a neighbour of which has already been deleted; so, if the repeat loop is executed i times, then i vertices have been introduced in set S'_{t+1} and, consequently, $n_{t+1} \leq \Delta + 1 + \sum_{j=1}^{i-1} \Delta$, or $i \geq (n_{t+1} - 1) / \Delta$.

By denoting by $\alpha'(G)$ the cardinality of the set S' finally provided by the algorithm, we get

$$\alpha'(G) = \sum_{j=1}^t \sigma_j + i \geq \sum_{j=1}^t \sigma_j + \frac{n_{t+1} - 1}{\Delta} \quad (2)$$

Let us recall that the variables of F for the vertices of G_{t+1} , are all equal to $1/2$, hence $v(F) = n_{t+1}/2$; moreover, since IS is a maximization problem, $\alpha(G_{t+1}) \leq v(F) = n_{t+1}/2$. On the other hand, if $\alpha(G_{t+1}) = v(F)$, then G_{t+1} is KE and the execution of the second if condition of Algorithm 1 would provide the optimal solution for G_{t+1} , the whole of the algorithm attaining an approximation ratio of 1 for IS; so, since an independent set in a graph is a set of discrete objects, $\alpha(G_{t+1}) \leq (n_{t+1}/2) - 1$ and, by using expression (1), we get

$$\alpha(G) \leq \sum_{j=1}^t \sigma_j + \frac{n_{t+1}}{2} - 1 \quad (3)$$

From expressions (1) and (3), we obtain

$$\frac{\alpha'(G)}{\alpha(G)} \geq \frac{\sum_{j=1}^t \sigma_j + \frac{n_{t+1} - 1}{\Delta}}{\sum_{j=1}^t \sigma_j + \frac{n_{t+1}}{2} - 1} \geq \frac{\frac{n_{t+1} - 1}{\Delta}}{\frac{n_{t+1}}{2} - 1} = \frac{2}{\Delta} \left(1 + \frac{1}{n_{t+1} - 2}\right) \geq \frac{2}{\Delta} \left(1 + \frac{1}{n - 2}\right) \quad (4)$$

Concerning the complexity of Algorithm 1, the call of the algorithm of [1] and the execution of the if condition take time $O(n^{2.5})$; the while loop will be executed at most n times, each time requiring $O(n^{2.5})$ [4]; finally, if the vertices are, a priori, sorted in increasing degree-order ($O(n \log n)$), the reorganization of the sorted list, after the deletion of the selected vertex and the update of the degrees, can be performed in $O(\log n)$ by using a heap (and this reorganization will be repeated, at most n times); on the other hand, the deletion of the edges takes, at most, $O(|E|)$; consequently, the whole of the complexity of the repeat loop takes $O(\max\{|E|, n \log n\})$. So, the total worst case complexity of Algorithm 1 is of $O(n^{3.5})$.

In fact, if the complexity of algorithm is $O(n^{3.5})$ (i.e., the while loop of Algorithm 1 is executed), the worst case approximation ratio of the algorithm is even better. Since the while loop is executed at least once, then at least a set S (corresponding to variables of F valued by 1) is formed; by supposing that G is connected, and by the constraints of IS, a set \bar{S} associated to S is also defined and at

least one vertex of $G[V_C \setminus (S \cup \bar{S})]$ has at least one neighbour in \bar{S} ; consequently, the minimum degree of $G[V_C \setminus (S \cup \bar{S})]$ is at most $\Delta - 1$ (moreover, the degrees of the survived vertices, for ulterior executions of the **while** loop of Algorithm 1, decrease). Hence, $i \geq n_{t+1} / \Delta$ (recall that i is the number of vertices added to S' during the **repeat** loop of Algorithm 1); on the other hand $\sum_{j=1}^t \sigma_j \geq 1$ and, since the function $(x+a)/(x+b)$, $b > a$ is increasing in x , expression (4) gives now $(\alpha'(G) / \alpha(G)) \geq (2 / \Delta) (1 + \Delta / n_{t+1})$. By the remark made above on the creation of the sets S and \bar{S} (resulting from the execution of the **while** loop of Algorithm 1), we have $n_{t+1} \leq n - 2$ (by supposing that there were only one execution of the **while** loop, this execution having produced sets S and \bar{S} , both of size 1); consequently

$$\frac{\alpha'(G)}{\alpha(G)} \geq \frac{2}{\Delta} \left(1 + \frac{\Delta}{n-2}\right)$$

Let us remark here that even if the improvement obtained in this note is slight and the quantity $2 / \Delta$ remains its lower bound, neither the fine analysis of [3], nor the analysis of [5], allow to obtain an improvement even as weak as the one obtained here.

We close this note by proving that for graphs admitting a perfect matching (of cardinality $\lfloor n/2 \rfloor$) the greedy algorithm (the **repeat** loop of Algorithm 1) attains approximation ratio strictly greater than $2 / \Delta$.

Let us consider a minimum vertex cover C^* and the corresponding maximum independent set S^* in a graph G . Let us also suppose that, given a matching M , there are f matching edges such that both their extremities belong to C^* , for the remaining ones, one of their extremities belonging to C^* and the other one to S^* . Let us call these edges "dissymmetric" and denote by F the set of theses "dissymmetric" edges ($|F| = f$). For M (in the case that is not perfect), let us denote by X the set of the exposed (non-saturated) vertices of G with respect to M , and by X_C ($|X_C| = g$) and X_S the subsets of X belonging to C^* and S^* , respectively (of course, $X = X_C \cup X_S$). The numbers f and g , consequently the sets F and X_C , depend not only on M but also, for a fixed matching, on the sets C^* (and S^*) considered. However, the sum $f + g$ is a quantity depending only on G ($f + g$ can be considered as the "discrete duality gap").

In fact, for every graph $|C^*| = \tau(G) = m + (f + g)$, $|S^*| = \alpha(G) = n - m - (f + g)$.

We have seen (see expression (2)) that the size of the solution delivered by the greedy IS-algorithm on a graph of order n is $\alpha'(G) \geq (n-1) / \Delta$. From the expression for $|S^*|$, we get $\alpha(G) = n - m - (f + g) = m + e - (f + g)$, where e is the number of the exposed vertices (of G) with respect to a maximum matching of G .

On the other hand, let us suppose that $f = n/x$; so, given that G admits a perfect matching M , we have the following for the terms of the third part of the last expression for $\alpha(G)$: $m \leq n/2$, $e - g \leq 1$ and $f = n/x$. From these last expressions, the last one for $\alpha(G)$, as well as the one for $\alpha'(G)$, we obtain $\alpha'(G) / \alpha(G) \geq \left[\frac{n-1}{\Delta} \right] / \left[\frac{n}{2} + 1 - \frac{n}{x} \right]$.

The function on the righthand side is decreasing in x ; on the other hand, we can, without loss of generality, suppose⁴ that $f = n/x \geq k$, for any fixed positive integer k , or $x \leq n/k$; so, after some easy algebra, we get

$$\frac{\alpha'(G)}{\alpha(G)} \geq \frac{2}{\Delta} \left(1 + \frac{2k+1}{n-2(k-1)}\right)$$

The second term in the parentheses can be quite large, always remaining, however, of $o(1)$. Moreover, this improvement for the approximation ratio of the greedy algorithm is obtained (if the information on the existence of a perfect matching is given) in time linear to the number of edges of G .

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⁴ In the opposite case, as we show in [2], IS is polynomial.

THE CONTINUOUS PROJECTION-GRADIENT METHOD OF THE FOURTH ORDER

Andjelija NEDIĆ

*Mathematical Faculty, University of Montenegro,
Cetinjski put, 81000 Podgorica, Yugoslavia*

Abstract:* In this paper the continuous projection-gradient method of the fourth order for solving the convex minimization problem in Euclidean space is considered. The sufficient conditions for convergence are established.

Keywords: Projection-gradient method, convexity.

1. INTRODUCTION

Consider the following minimization problem

$$J(u) \rightarrow \inf, \quad u \in U, \quad (1)$$

where U is a closed, convex subset of a real Euclidean space \mathbb{E}^n , function $J(u)$ is continuously differentiable and convex on \mathbb{E}^n . The scalar product of two elements $u, v \in \mathbb{E}^n$ will be denoted by: $\langle u, v \rangle$; $\|u\| = \langle u, u \rangle^{1/2}$ is the norm of the element u . Suppose that

$$J_* = \inf_{u \in U} J(u) > -\infty, \quad U_* = \{u \in U : J(u) = J_*\} \neq \emptyset. \quad (2)$$

The continuous minimization methods of the projection-gradient type

$$u'(t) + u(t) = P_U [u(t) - \alpha(t) J'(u(t))], \quad t \geq 0,$$

$$u(0) = u_0, \quad u_0 \in \mathbb{E}^n,$$

have been proposed and investigated in [5, 6] for $U = \mathbb{E}^n$ and in [1] for $U \subseteq \mathbb{E}^n$, $\alpha(t) = \alpha > 0$, $t \geq 0$. The further investigation in this area, considering the continuous projection-gradient methods of the second and third order has been presented in [1, 2]. This paper presents the continuous projection-gradient method of the fourth order.

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2. THE CONDITIONS FOR CONVERGENCE

For solving the problem (1) we will use the continuous projection gradient method of the fourth order

$$\begin{aligned} & \beta_4(t) u^{iv} + \beta_3(t) u''' + \beta_2(t) u'' + u' + u = \\ & P_U [u - \alpha(t) J'(u)], \quad t \geq 0, \\ & u(0) = u_0, \quad u'(0) = u_1, \quad u''(0) = u_2, \quad u'''(0) = u_3 \end{aligned} \quad (3)$$

where $P_U(z)$ – denotes the projection of the point z on the set U ; $u_i, i = 0, 1, 2, 3$ are given initial points from the Euclidean space \mathbb{E}^n ; $\alpha(t), \beta_i(t), i = 2, 3, 4$ are the parameters of the method (3), $u = u(t), u^{(i)}(t) = d^i u(t) / dt^i, i = 1, 2, 3, 4, J'(u)$ – gradient of the function $J(u)$.

THEOREM 1. Suppose that

- 1) U is a convex closed set in Euclidean space \mathbb{E}^n ; function $J(u)$ is convex and differentiable on \mathbb{E}^n ; the gradient $J'(u)$ satisfies the Lipschitz condition

$$\|J'(u) - J'(v)\| \leq L \|u - v\|, \quad u, v \in \mathbb{E}^n, \quad L = \text{const} < +\infty;$$

the conditions (2) are satisfied;

- 2) the parameters $\alpha(t), \beta_i(t), i = 2, 3, 4$ of the method (3) are such that

$$\begin{aligned} & \alpha(t) \in \mathbb{C} [0, +\infty), 0 < \alpha_0 \leq \alpha(t) \leq \alpha_1, \quad t \geq 0; \\ & \beta_2(t) \in \mathbb{C}^2 [0, +\infty), \quad \beta_3(t) \in \mathbb{C}^3 [0, +\infty), \quad \beta_4(t) \in \mathbb{C}^4 [0, +\infty), \\ & \beta_i'(t) \leq 0, \quad \beta_i''(t) \geq 0, \quad i = 2, 3, 4, \quad t \geq 0; \\ & \beta_i'''(t) \leq 0, \quad i = 3, 4; \quad \beta_i^{iv}(t) \geq 0, \quad t \geq 0; \\ & \lim_{t \rightarrow \infty} \beta_i(t) = \beta_{i\infty} > 0, \quad i = 2, 3, 4; \end{aligned}$$

$$1 - \alpha_1 L - \beta_{2\infty} > 0, \quad \beta_{2\infty}^2 (1 - \alpha_1 L) + \beta_{4\infty} - 2\beta_{3\infty} > 0,$$

$$\beta_{3\infty}^2 (1 - \alpha_1 L) + 2\beta_{2\infty} \beta_{4\infty} > 0, \quad \beta_{2\infty} - \frac{3}{2} \beta_{3\infty} - \beta_{3\infty} \beta_{4\infty} > 0,$$

$$\beta_{2\infty} \beta_{3\infty} - 3\beta_{4\infty} - \beta_{2\infty} \beta_{4\infty}^2 > 0.$$

Then for any initial values $u_i \in \mathbb{E}^n, i = 0, 1, 2, 3$ there is a point $u_\infty \in U$ such that

$$\begin{aligned} & \lim_{t \rightarrow \infty} \left\{ \sum_{i=1}^4 \|u^{(i)}(t)\| + \|u(t) - u_\infty\| \right\} = 0, \\ & \int_0^{+\infty} \left\{ \sum_{i=1}^4 \|u^{(i)}(s)\|^2 + f(s) \|u(s) - u_\infty\|^2 \right\} ds < +\infty \end{aligned}$$

where $f(s) = \beta_2''(s) - \beta_3'''(s) + \beta_4^{iv}(s)$, for all $s \geq 0$.

PROOF. Note that there exist functions, for example

$$\alpha(t) = \alpha_0 \frac{2+t}{1+t}, \quad \beta_i(t) = \beta_{i\infty} + \frac{1}{1+t}, \quad \alpha_0 > 0, \quad \beta_{i\infty} > 0,$$

for $i = 2, 3, 4$; and $t \geq 0$, which satisfy the conditions of the theorem.

From the inequalities for the derivatives of $\beta_i(t)$ and $\beta_{i\infty} > 0$, $i = 2, 3, 4$; it can be proved that

$$\begin{aligned} \lim_{t \rightarrow \infty} \beta_i'(t) &= 0, \quad i = 2, 3, 4 & \lim_{t \rightarrow \infty} \beta_i''(t) &= 0, \quad i = 2, 3; \\ \lim_{t \rightarrow \infty} \beta_4'''(t) &= 0. \end{aligned} \quad (4)$$

Besides that, from the conditions on the limits $\beta_{i\infty}$ and α_1 we can find

$$\begin{aligned} 1 > \beta_{2\infty} > \frac{3}{2} \beta_{3\infty} > \frac{9}{2} \beta_{4\infty}, & \quad \beta_{3\infty} - 3 \beta_{4\infty}^2 > 0, \\ 1 - \beta_{2\infty} \beta_{4\infty} - \beta_{4\infty}^4 > 0, & \quad \beta_{3\infty} - 2 \beta_{4\infty} - 2 \beta_{4\infty}^2 > 0, \\ \beta_{2\infty} - 2 \beta_{3\infty} \beta_{4\infty} > 0, & \quad \beta_{3\infty} - \beta_{2\infty}^2 \beta_{4\infty} - \beta_{4\infty} > 0. \end{aligned} \quad (5)$$

As we know [3], there is the unique solution $u = u(t)$, $t \geq 0$ to the differential equation (3) for any given initial values $u_i \in \mathbb{E}^n$, $i = 0, 1, 2, 3$. For every $u_* \in U_*$, it holds: (see [4], pp.165, Theorem 3)

$$\langle J'(u_*), w - u_* \rangle \geq 0, \quad w \in U. \quad (6)$$

From (3) and the property of the projection operator (see [4]) it can be derived

$$\begin{aligned} &\langle \beta_4(t) u^{iv} + \beta_3(t) u''' + \beta_2(t) u'' + u' + \alpha(t) J'(u), \\ &\beta_4(t) u^{iv} + \beta_3(t) u''' + \beta_2(t) u'' + u' + u - v \rangle \leq 0, \\ &v \in U, \quad t \geq 0. \end{aligned} \quad (7)$$

Multiplying the inequality (6) by $[-\alpha(t)]$ for $w = \beta_4(t) u^{iv} + \beta_3(t) u''' + \beta_2(t) u'' + u' + u$ and summing it with (7) for $v = u_*$, we have

$$\begin{aligned} &\langle \beta_4(t) u^{iv} + \beta_3(t) u''' + \beta_2(t) u'' + u' + \alpha(t) [J'(u) - J'(u_*)], \\ &\beta_4(t) u^{iv} + \beta_3(t) u''' + \beta_2(t) u'' + u' + u - u_* \rangle \leq 0, \\ &t \geq 0, \quad u_* \in U_*. \end{aligned}$$

The last inequality can be written in the following way

$$\begin{aligned} &\beta_4^2(t) \|u^{iv}\|^2 + 2\beta_4(t)\beta_3(t) \langle u^{iv}, u''' \rangle + 2\beta_4(t)\beta_2(t) \langle u^{iv}, u'' \rangle + \\ &2\beta_4(t) \langle u^{iv}, u' \rangle + \beta_4(t) \langle u^{iv}, u - u_* \rangle + \beta_3^2(t) \|u'''\|^2 + \\ &2\beta_3(t)\beta_2(t) \langle u''', u'' \rangle + 2\beta_3(t) \langle u''', u' \rangle + \\ &\beta_3(t) \langle u''', u - u_* \rangle + \beta_2^2(t) \|u''\|^2 + 2\beta_2(t) \langle u'', u' \rangle + \\ &\beta_2(t) \langle u'', u - u_* \rangle + \|u'\|^2 + \langle u', u - u_* \rangle \leq \\ &\alpha(t) \langle J'(u) - J'(u_*), u_* - [u + u' + \beta_2(t)u'' + \beta_3(t)u''' + \beta_4(t)u^{iv}] \rangle \\ &t \geq 0, \quad u_* \in U_*. \end{aligned} \quad (8)$$

Since the function $J(u)$ is convex, differentiable and its gradient $J'(u)$ satisfies the Lipschitz's condition we have (see [4], pp.175).

$$\langle J'(u) - J'(v), v - w \rangle \leq \frac{L}{4} \|u - w\|^2, \quad u, v, w \in \mathbb{E}^n \quad (9)$$

In order to write (8) in a more compact way we will use the denotations

$$\begin{aligned} r(t) &= \|u'''(t)\|^2, \\ q(t) &= \|u''(t)\|^2, \\ p(t) &= \|u'(t)\|^2, \\ x(t, u_*) &= \frac{1}{2} \|u(t) - u_*\|^2. \end{aligned} \quad (10)$$

Using (10), (9) and inequality $(a + b + c + d)^2 \leq 4(a^2 + b^2 + c^2 + d^2)$, which is true for all $a, b, c, d \in \mathbb{R}$, from (8) we will get

$$\begin{aligned} &\beta_4^2(t) [1 - \alpha(t)L] \|u^{iv}\|^2 + \beta_4(t) \beta_3(t) r'(t) + \\ &[\beta_3^2(t) (1 - \alpha(t)L) - 2\beta_4(t) \beta_2(t)] r(t) + \beta_4(t) \beta_2(t) q''(t) + \\ &[\beta_3(t) \beta_2(t) - 3\beta_4(t)] q'(t) + [\beta_2^2(t) (1 - \alpha(t)L) + \beta_4(t) - 2\beta_3(t)] q(t) + \\ &\beta_4(t) p'''(t) + [\beta_3(t) - 2\beta_4(t)] p''(t) + [\beta_2(t) - \frac{3}{2}\beta_3(t)] p'(t) + \\ &[1 - \alpha(t)L - \beta_2(t)] p(t) + \\ &\beta_4(t) x^{iv}(t, u_*) + \beta_3(t) x'''(t, u_*) + \beta_2(t) x''(t, u_*) + x'(t, u_*) \leq 0, \\ &t \geq 0, \quad u_* \in U_* \end{aligned}$$

After the integration on the segment $[\xi, t]$, $t > \xi$, where $\xi \geq 0$ is arbitrary, taking into account the conditions 2) of the Theorem, we have

$$\begin{aligned} &\int_{\xi}^t \{ \beta_4^2(s) [1 - \alpha_1 L] \|u^{iv}(s)\|^2 + [\beta_3^2(s) (1 - \alpha_1 L) - 2\beta_4(s) \beta_2(s)] r(s) + \\ &[\beta_2^2(s) (1 - \alpha_1 L) + \beta_4(s) - 2\beta_3(s) + 3\beta_4'(s)] q(s) + \\ &[1 - \alpha_1 L - \beta_2(s) + \frac{3}{2}\beta_3'(s) - 2\beta_4''(s)] p(s) + \\ &[\beta_2''(s) - \beta_3''(s) + \beta_4^{iv}(s)] x(s, u_*) \} ds + \\ &\beta_4(t) \beta_3(t) r(t) + \beta_4(t) \beta_2(t) q'(t) + \\ &[\beta_3(t) \beta_2(t) - 3\beta_4(t)] q(t) + \beta_4(t) p''(t) + \\ &[\beta_3(t) - 2\beta_4(t) - \beta_4'(t)] p'(t) + [\beta_2(t) - \frac{3}{2}\beta_3(t) + 2\beta_4'(t)] p(t) + \\ &\beta_4(t) x'''(t, u_*) + [\beta_3(t) - \beta_4'(t)] x''(t, u_*) + [\beta_2(t) - \beta_3'(t) + \beta_4''(t)] x'(t, u_*) + \\ &x(t, u_*) \leq C_0(\xi, u_*), \\ &t > \xi \geq 0, \quad u_* \in U_* \end{aligned} \quad (11)$$

where

$$\begin{aligned}
C_0(\xi, u_*) &= \beta_4(\xi) \beta_3(\xi) r(\xi) + \beta_4(\xi) \beta_2(\xi) q'(\xi) + \\
&\quad [\beta_3(\xi) \beta_2(\xi) - 3\beta_4(\xi) - (\beta_2(\xi) \beta_4(\xi))'] q(\xi) + \beta_4(\xi) p''(\xi) + \\
&\quad [\beta_3(\xi) - 2\beta_4(\xi) - \beta_4'(\xi)] p'(\xi) + \\
&\quad [\beta_2(\xi) - \frac{3}{2}\beta_3(\xi) + 2\beta_4'(\xi) - \beta_3'(\xi) + \beta_4''(\xi)] p(\xi) + \\
&\quad \beta_4(\xi) x''(\xi, u_*) + [\beta_3(\xi) - \beta_4'(\xi)] x'(\xi, u_*) + \\
&\quad [\beta_2(\xi) - \beta_3'(\xi) + \beta_4''(\xi)] x(\xi, u_*) + \\
&\quad [1 - \beta_2'(\xi) + \beta_3''(\xi) - \beta_4'''(\xi)] x(\xi, u_*).
\end{aligned} \tag{12}$$

From the condition 2) of the Theorem and (4), (5), it follows

$$\begin{aligned}
\beta_4^2(s) [1 - \alpha_1 L] &> 0, \quad s \geq 0, \\
\lim_{s \rightarrow \infty} [\beta_3^2(s) (1 - \alpha_1 L) - 2\beta_4(s) \beta_2(s)] &= \\
\beta_{3\infty}^2 (1 - \alpha_1 L) - 2\beta_{4\infty} \beta_{2\infty} &> 0, \\
\lim_{s \rightarrow \infty} [\beta_2^2(s) (1 - \alpha_1 L) + \beta_4(s) - 2\beta_3(s) + 3\beta_4'(s)] &= \\
\beta_{2\infty}^2 (1 - \alpha_1 L) + \beta_{4\infty} - 2\beta_{3\infty} &> 0, \\
\lim_{s \rightarrow \infty} [1 - \alpha_1 L - \beta_2(s) + \frac{3}{2}\beta_3'(s) - 2\beta_4''(s)] &= \\
1 - \alpha_1 L - \beta_{2\infty} &> 0.
\end{aligned}$$

Therefore there are ε , $0 < \varepsilon < 1/2$ and $t_0 \geq 0$, such that:

$$\begin{aligned}
\beta_4^2(s) [1 - \alpha_1 L] &\geq \varepsilon, \\
\beta_3^2(s) (1 - \alpha_1 L) - 2\beta_4(s) \beta_2(s) &\geq \varepsilon, \\
\beta_2^2(s) (1 - \alpha_1 L) + \beta_4(s) - 2\beta_3(s) + 3\beta_4'(s) &\geq \varepsilon, \\
1 - \alpha_1 L - \beta_2(s) + \frac{3}{2}\beta_3'(s) - 2\beta_4''(s) &\geq \varepsilon, \\
s &\geq t_0.
\end{aligned}$$

Then from (11) we have

$$\begin{aligned}
\varepsilon \int_{\xi}^t \left\{ \sum_{i=1}^4 \|u^{(i)}(s)\|^2 + f(s) \|u(s) - u_*\|^2 \right\} ds + \\
\beta_4(t) \beta_3(t) r(t) + \beta_4(t) \beta_2(t) q'(t) + \\
[\beta_3(t) \beta_2(t) - 3\beta_4(t)] q(t) + \beta_4(t) p''(t) + \\
[\beta_3(t) - 2\beta_4(t) - \beta_4'(t)] p'(t) + [\beta_2(t) - \frac{3}{2}\beta_3(t) + 2\beta_4'(t)] p(t) + \\
\beta_4(t) x''(t, u_*) + [\beta_3(t) - \beta_4'(t)] x'(t, u_*) + [\beta_2(t) - \beta_3'(t) + \beta_4''(t)] x(t, u_*) + \\
x(t, u_*) \leq C_0(\xi, u_*),
\end{aligned} \tag{13}$$

for all $t > \xi \geq 0$, $u_* \in U_*$. Let $h(t) = \exp \left\{ \int_0^t \beta_4(s) ds \right\} > 0$, $t \geq 0$. In (13) the integral is positive and $\beta_4(t) \beta_3(t) r(t) \geq 0$, $t \geq 0$, so they can be omitted. After that we will multiply (13) by $h(t)$ and integrate it on the segment $[\xi, t]$, for $t > \xi \geq t_0$. In this way we will get

$$\begin{aligned} & \beta_4(t) \beta_2(t) h(t) q(t) + \\ & \int_{\xi}^t [\beta_3(s) \beta_2(s) - 3 \beta_4(s) - (\beta_2(s) \beta_4(s))' - \beta_2(s) \beta_4^2(s)] h(s) q(s) ds + \\ & \beta_4(t) h(t) p'(t) + [\beta_3(t) - 2 \beta_4(t) - 2 \beta_4'(t) - \beta_4^2(t)] h(t) p(t) + \\ & \int_{\xi}^t \{ [\beta_4(s) h(s)]'' - [(\beta_3(s) - 2 \beta_4(s) - \beta_4'(s)) h(s)]' + \\ & \quad [\beta_2(s) - \frac{3}{2} \beta_3(s) + 2 \beta_4'(s)] h(s) \} p(s) ds + \\ & \beta_4(t) h(t) x'''(t, u_*) + [\beta_3(t) - 2 \beta_4'(t) - \beta_4^2(t)] h(t) x'(t, u_*) + \\ & \{ [\beta_2(t) - \beta_3'(t) + \beta_4''(t)] h(t) - [(\beta_3(t) - \beta_4'(t)) h(t)]' + [\beta_4(t) h(t)]'' \} x(t, u_*) + \\ & \int_{\xi}^t \{ h(s) - [(\beta_2(s) - \beta_3'(s) + \beta_4''(s)) h(s)]' + \\ & \quad [(\beta_3(s) - \beta_4'(s)) h(s)]'' - [\beta_4(s) h(s)]''' \} x(s, u_*) ds \leq \\ & C_0(\xi, u_*) \int_{\xi}^t h(s) ds + C_1(\xi, u_*), \\ & t > \xi \geq t_0, \quad u_* \in U_* \end{aligned}$$

where

$$\begin{aligned} C_1(\xi, u_*) = & \beta_4(\xi) \beta_2(\xi) h(\xi) q(\xi) + \beta_4(\xi) h(\xi) p'(\xi) + \\ & [\beta_3(\xi) - 2 \beta_4(\xi) - 2 \beta_4'(\xi) - \beta_4^2(\xi)] h(\xi) p(\xi) + \\ & \beta_4(\xi) h(\xi) x'''(\xi, u_*) + [\beta_3(\xi) - 2 \beta_4'(\xi) - \beta_4^2(\xi)] h(\xi) x'(\xi, u_*) + \\ & \{ [\beta_2(\xi) - \beta_3'(\xi) + \beta_4''(\xi)] h(\xi) - [(\beta_3(\xi) - \beta_4'(\xi)) h(\xi)]' + \\ & \quad [\beta_4(\xi) h(\xi)]'' \} x(\xi, u_*). \end{aligned}$$

From the conditions 2) of the Theorem, (4) and (5) it can be shown that all integrals on the left hand side of the above inequality are non-negative for some $t_1 \geq t_0$ and every $t > \xi \geq t_1$. Besides that we have: $\beta_4'(t) \leq 0$, $\beta_3'(t) \leq 0$, $\beta_4''(t) \geq 0$, so the last inequality becomes

$$\begin{aligned}
& \beta_4(t) \beta_2(t) h(t) q(t) + \beta_4(t) h(t) p'(t) + \\
& [\beta_3(t) - 2\beta_4(t) - \beta_4^2(t)] h(t) p(t) + \\
& \beta_4(t) h(t) x''(t, u_*) + [\beta_3(t) - 2\beta_4(t) - \beta_4^2(t)] h(t) x'(t, u_*) + \\
& [\beta_2(t) - \beta_3(t) \beta_4(t) + 4\beta_4(t) \beta_4'(t)] h(t) x(t, u_*) \leq \\
& C_0(\xi, u_*) \int_{\xi}^t h(s) ds + C_1(\xi, u_*),
\end{aligned} \tag{14}$$

for every $t > \xi \geq t_1, u_* \in U_*$. Integrating (14) on $[\xi, t]$, we have

$$\begin{aligned}
& \int_{\xi}^t \beta_4(s) \beta_2(s) h(s) q(s) ds + \int_{\xi}^t [\beta_3(s) - 2\beta_4(s) - 2\beta_4^2(s) - \beta_4'(s)] h(s) p(s) ds + \\
& \beta_4(t) h(t) p(t) + \beta_4(t) h(t) x'(t, u_*) + [\beta_3(t) - 3\beta_4'(t) - 2\beta_4^2(t)] h(t) x(t, u_*) + \\
& \int_{\xi}^t \{ [\beta_2(s) - \beta_3(s) \beta_4(s) + 4\beta_4(s) \beta_4'(s)] h(s) - \\
& [(\beta_3(s) - 2\beta_4'(s) - \beta_4^2(s)) h(s)]' + [\beta_4(s) h(s)]'' \} x(s, u_*) ds \leq \\
& C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s h(\theta) d\theta ds + C_1(\xi, u_*) (t - \xi) + C_2(\xi, u_*),
\end{aligned}$$

for $t > \xi \geq t_1, u_* \in U_*$ where

$$\begin{aligned}
C_2(\xi, u_*) = & \beta_4(\xi) h(\xi) p(\xi) + \beta_4(\xi) h(\xi) x'(\xi, u_*) + \\
& [\beta_3(\xi) - 3\beta_4'(\xi) - 2\beta_4^2(\xi)] h(\xi) x(\xi, u_*).
\end{aligned}$$

Taking into account the conditions 2) of the Theorem and relations (4), (5), we can find that there exists $t_2 \geq t_1$, such that the integrals on the left hand side of the last inequality are non-negative for all $t > \xi \geq t_2$. Hence

$$\begin{aligned}
& \beta_4(t) h(t) p(t) + \beta_4(t) h(t) x'(t, u_*) + \\
& [\beta_3(t) - 3\beta_4'(t) - 2\beta_4^2(t)] h(t) x(t, u_*) \leq \\
& (15) \quad C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s h(\theta) d\theta ds + C_1(\xi, u_*) + (t - \xi) + C_2(\xi, u_*),
\end{aligned}$$

$$t > \xi \geq t_2, \quad u_* \in U_*.$$

After the integration in (15) on $[\xi, t]$, we get

$$\begin{aligned}
& \int_{\xi}^t \{ \beta_4(s) p(s) + [\beta_3(s) - 4\beta_4'(s) - 3\beta_4^2(s)] x(t, u_*) \} h(s) ds + \beta_4(t) h(t) x(t, u_*) \leq \\
& C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s \int_{\xi}^v h(\theta) d\theta dv ds + \frac{1}{2} C_1(\xi, u_*) (t - \xi)^2 + C_2(\xi, u_*) (t - \xi) + C_3(\xi, u_*),
\end{aligned}$$

for $t > \xi \geq t_2, u_* \in U_*$, where

$$C_3 = \beta_4(\xi) h(\xi) x(\xi, u_*).$$

From $\beta_4' \leq 0$ and (5) it follows that there is $t_3, t_3 \geq t_2$, such that the integrals on the left hand side of the above inequality are non-negative for all $t > \xi \geq t_3$, so that

$$\beta_4(t) h(t) x(t, u_*) \leq C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s \int_{\xi}^v h(\theta) d\theta dv ds + \\ \frac{1}{2} C_1(\xi, u_*) (t - \xi)^2 + C_2(\xi, u_*) (t - \xi) + C_3(\xi, u_*),$$

for $t > \xi \geq t_3, u_* \in U_*$. Consequently

$$\overline{\lim}_{t \rightarrow \infty} \|u(t) - u_*\|^2 \leq 2 \overline{\lim}_{t \rightarrow \infty} [\beta_4(t) h(t)]^{-1} \{ C_3(\xi, u_*) + C_2(\xi, u_*) (t - \xi) + \\ \frac{1}{2} C_1(\xi, u_*) (t - \xi)^2 + C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s \int_{\xi}^v h(\theta) d\theta dv ds \} = \\ 2 \beta_{4\infty}^{-1} \lim_{t \rightarrow \infty} h^{-1}(t) \{ C_3(\xi, u_*) + C_2(\xi, u_*) (t - \xi) + \\ \frac{1}{2} C_1(\xi, u_*) (t - \xi)^2 + C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s \int_{\xi}^v h(\theta) d\theta dv ds \}.$$

It is clear that $\lim_{t \rightarrow \infty} h(t) = +\infty$. In order to calculate the limit of the right hand side we will use the L'Hospital rule three times. After that we will get

$$\overline{\lim}_{t \rightarrow \infty} \|u(t) - u_*\|^2 \leq b_0 C_0(\xi, u_*), \quad (16) \\ \xi \geq t_3, \quad u_* \in U_*.$$

where $b_0 = 2 \beta_{4\infty}^{-4}$. Using (5) and $\beta_4' \leq 0$, it is not difficult to see that the third term on the left hand side of (15) is non-negative for some $t_4 \geq t_3$ and all $t \geq t_4$. From here and denotations (10) we have

$$\|u'(t)\|^2 + \langle u'(t), u(t) - u_* \rangle \leq \\ 2 \{ C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s h(\theta) d\theta ds + C_1(\xi, u_*) (t - \xi) + C_2(\xi, u_*) \} [\beta_4(t) h(t)]^{-1}, \\ t > \xi \geq t_4, \quad u_* \in U_*.$$

Therefore from (16) and the inequality

$$2 |\langle a, b \rangle| \leq \|a\|^2 + \|b\|^2, \quad a, b \in \mathbf{E}^n, \quad (17)$$

it follows

$$\begin{aligned} \overline{\lim}_{t \rightarrow \infty} \|u'(t)\|^2 &\leq b_0 C_0(\xi, u_*) + \\ 2 \overline{\lim}_{t \rightarrow \infty} [\beta_4(t) h(t)]^{-1} \{ &C_2(\xi, u_*) + C_1(\xi, u_*) (t - \xi) + C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s h(\theta) d\theta ds \} = \\ b_0 C_0(\xi, u_*) + \\ \frac{2}{\beta_{4\infty}} \overline{\lim}_{t \rightarrow \infty} \{ &C_0(\xi, u_*) \int_{\xi}^t \int_{\xi}^s h(\theta) d\theta ds + C_1(\xi, u_*) (t - \xi) + C_2(\xi, u_*) \} h^{-1}(t). \end{aligned}$$

Using the L'Hospital rule two times, it can be obtained

$$\begin{aligned} \overline{\lim}_{t \rightarrow \infty} \|u'(t)\|^2 &\leq b_1 C_0(\xi, u_*) \\ \xi &\geq t_4, \quad u_* \in U_*, \end{aligned} \tag{18}$$

where $b_1 = b_0 + 2/\beta_{4\infty}^3$. From (4), (5) it is clear that there is $t_5 \geq t_4$, such that the third and the sixth terms on the left hand side of (14) are non-negative. Taking into account (10) and (17) in a similar way as before, we can get

$$\begin{aligned} \overline{\lim}_{t \rightarrow \infty} \beta_4(t) \left[\beta_2(t) - \frac{3}{2} \beta_4(t) \right] \|u''(t)\|^2 &\leq \\ \overline{\lim}_{t \rightarrow \infty} \left\{ \frac{1}{2} [1 + |f_1(t)|] \|u - u_*\|^2 + \left[1 + \frac{|f_1(t)|}{2} \right] \|u'(t)\|^2 + \right. \\ \left. [h(t)]^{-1} \left[C_0(\xi, u_*) \int_{\xi}^t h(s) ds + C_1(\xi, u_*) \right] \right\}, \end{aligned}$$

where $f_1(t) = \beta_3(t) - 2\beta_4'(t) - \beta_4^2(t)$, $t \geq 0$. Let $M_1 > 0$, such that $|f_1(t)| \leq M_1$, $t \geq 0$. Then

$$\begin{aligned} \beta_{4\infty} \left[\beta_{2\infty} - \frac{3}{2} \beta_{4\infty} \right] \overline{\lim}_{t \rightarrow \infty} \|u''(t)\|^2 &\leq \\ \overline{\lim}_{t \rightarrow \infty} \left\{ \frac{1 + M_1}{2} \|u - u_*\|^2 + \left[1 + \frac{M_1}{2} \right] \|u'(t)\|^2 + \right. \\ \left. [h(t)]^{-1} \left[C_0(\xi, u_*) \int_{\xi}^t h(s) ds + C_1(\xi, u_*) \right] \right\}. \end{aligned}$$

From (5) we have $\beta_{4\infty} [\beta_{2\infty} - (3/2)\beta_{4\infty}] > 0$. Using L'Hospital's rule on the right hand side of the above inequality and the estimates (16), (18) we will get

$$\begin{aligned} \overline{\lim}_{t \rightarrow \infty} \|u''(t)\|^2 &\leq b_2 C_0(\xi, u_*), \\ \xi &\geq t_5, \quad u_* \in U_*, \end{aligned} \tag{19}$$

where b_2 does not depend on ξ and u_* . Taking into account (4), (5), (17) and denotations (10), from (13) we can obtain

$$\begin{aligned} & \varepsilon \int_{\xi}^t \left\{ \sum_{i=1}^4 \|u^{(i)}(s)\|^2 + f(s) \|u(s) - u_{\infty}\|^2 \right\} ds + \\ & [\beta_4(t)\beta_3(t) - \beta_4^2(t)\beta_2^2(t) - \beta_4^2(t)] \|u'''\|^2 \leq \\ & |f_2(t)| \|u''\|^2 + |f_3(t)| \|u'\|^2 + |f_4(t)| \|u - u_{\infty}\|^2 + C_0(\xi, u_{\infty}), \\ & t > \xi \geq t_5, \quad u_{\infty} \in U_{\infty}, \end{aligned}$$

where the functions f_i , $i = 2, 3, 4$; are bounded and do not depend on ξ and u_{∞} . Let $M_2 > 0$, such that $|f_i(t)| \leq M_2$, $i = 2, 3, 4$; $t \geq 0$. Since

$$\lim_{t \rightarrow \infty} [\beta_3(t) - \beta_4(t)\beta_2^2(t) - \beta_4^2(t)] = \beta_{3\infty} - \beta_{4\infty}\beta_{2\infty}^2 - \beta_{4\infty}^2 > 0,$$

then there exist a moment $t_7 \geq t_6$ and a number δ , $0 < \delta < \varepsilon$, such that for all $t \geq t_7$: $\beta_4(t)[\beta_3(t) - \beta_4(t)\beta_2^2(t) - \beta_4^2(t)] > \delta$. Hence

$$\begin{aligned} & \delta \left\{ \int_{\xi}^t \left[\sum_{i=1}^4 \|u^{(i)}(s)\|^2 + f(s) \|u(s) - u_{\infty}\|^2 \right] ds + \|u'''\|^2 \right\} \leq \\ & M_2 [\|u''\|^2 + \|u'\|^2 + \|u - u_{\infty}\|^2] + C_0(\xi, u_{\infty}), \end{aligned}$$

for all $t > \xi \geq t_7$, $u_{\infty} \in U_{\infty}$. Consequently

$$\delta \left\{ \int_{\xi}^t \left[\sum_{i=1}^4 \|u^{(i)}(s)\|^2 + f(s) \|u(s) - u_{\infty}\|^2 \right] ds \right\} \leq b_3 C_0(\xi, u_{\infty}), \quad (20)$$

$$\overline{\lim}_{t \rightarrow \infty} \|u'''\|^2 \leq b_3 C_0(\xi, u_{\infty}), \quad (21)$$

$$\xi \geq t_7, \quad u_{\infty} \in U_{\infty}.$$

where $b_3 = \delta^{-1} [M_2(b_2 + b_1 + b_0) + 1]$. From (20) it follows that there is a sequence $\{s_j\} \subseteq [0, +\infty)$ such that

$$\lim_{s \rightarrow \infty} \left[\sum_{i=1}^4 \|u^{(i)}(s)\|^2 \right] = \lim_{j \rightarrow \infty} \left[\sum_{i=1}^4 \|u^{(i)}(s_j)\|^2 \right] = 0,$$

i.e.

$$\lim_{j \rightarrow \infty} \|u^{(i)}(s_j)\| = 0, \quad i = 1, 2, 3, 4. \quad (22)$$

Then from (16) and the conditions 2) of the Theorem, it is obvious that there exist a subsequence $\{\hat{s}_j\}$, a point $u_{\infty} \in \mathbb{E}^n$ and a real number $\alpha_{\infty} > 0$, such that

$$\lim_{j \rightarrow \infty} \|u(\hat{s}_j) - u_{\infty}\| = 0, \quad \lim_{j \rightarrow \infty} \alpha(\hat{s}_j) = \alpha_{\infty}. \quad (23)$$

Setting $t = \hat{s}_j$ in the differential equation (3), from (22) and (23), we can get

$$\begin{aligned} & \lim_{j \rightarrow \infty} \|\beta_4(\hat{s}_j) u^{iv}(\hat{s}_j) + \beta_3(\hat{s}_j) u'''\!(\hat{s}_j) + \beta_2(\hat{s}_j) u''(\hat{s}_j) + u'(\hat{s}_j) + u(\hat{s}_j) - \\ & P_U [u(\hat{s}_j) - \alpha(\hat{s}_j) J'(u(\hat{s}_j))]\| = \\ & \|u_{\infty} - P_U [u_{\infty} - \alpha_{\infty} J'(u_{\infty})]\| = 0. \end{aligned}$$

Consequently (see [4], pp.171) $u_\infty \in U_*$. From (12) where $\xi = \hat{s}_j$, $u_* = u_\infty \in U_*$, along with (22), (23) we have

$$\lim_{j \rightarrow \infty} C_0(\hat{s}_j, u_\infty) = 0.$$

Let $j_0 \geq 1$ be such a number that $\hat{s}_j \geq t_7$, for every $j \geq j_0$. Then the relations (16), (18)-(21) hold for $\xi = \hat{s}_j$, $j \geq j_0$ and $u_* = u_\infty$. Therefore

$$\overline{\lim}_{t \rightarrow \infty} \|u(t) - u_\infty\|^2 \leq b_0 \lim_{j \rightarrow \infty} C_0(\hat{s}_j, u_\infty) = 0.$$

$$\overline{\lim}_{t \rightarrow \infty} \|u^{(i)}(t)\|^2 \leq b_i \lim_{j \rightarrow \infty} C_0(\hat{s}_j, u_\infty) = 0, \quad i = 1, 2, 3.$$

Hence

$$\lim_{t \rightarrow \infty} \|u(t) - u_\infty\| = 0, \quad \lim_{t \rightarrow \infty} \|u^{(i)}(t)\| = 0, \quad i = 1, 2, 3. \quad (24)$$

From the differential equation (3) and the relations, when $t \rightarrow \infty$, it can be obtained

$$\lim_{t \rightarrow \infty} \|u^{iv}(t)\| = 0. \quad (25)$$

The inequality (20) and the relations (24), (25) give the statement of the theorem.

3. CONCLUSION

The projection-gradient methods of the higher order are important because of their higher rate of convergence in comparison with the methods of that type of the first order [1]. Besides that, the continuous methods give us a large choice of the numerical integration methods for solving the corresponding differential equations. These facts justify the investigation of the continuous methods of the higher order.

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