

$$A_2 = \begin{bmatrix} 1 & \dots & 1 & \dots & 1 \\ 1 & \dots & 1 & \dots & 1 \\ 0 & \dots & \dots & 0 & \dots \\ 1 & \dots & 1 & \dots & 1 \end{bmatrix}$$

Let us denote by

$$V_1 = \{x \in R^{mn} \mid A_1 x = a\}$$

$$V_2 = \{x \in R^{mn} \mid A_2 x = b\}$$

$$V = V_1 \cap V_2$$

$$X = \{x \in R^{mn} \mid x_{ij} \geq 0\}.$$

Then, the feasible set is

$$D = V \cap X.$$

In this paper it is shown that projection point onto V could be reached in 2 steps after consecutive projections onto V_1 and V_2 , using simple formulas (with complexities $O(m)$ and $O(n)$).

2. PROJECTION OF A POINT ONTO A TRANSPORTATION MANIFOLD

In this section we shall consider the problem

$$(P_V) \quad \inf \frac{1}{2} \|x - y\|^2, \quad x \in V,$$

i.e.

$$(P_{V_j}) \quad \inf \frac{1}{2} \|x - y\|^2, \quad x \in V_j, \quad j = 1, 2.$$

where y is any point in R^{mn} . Optimal solution of (P_V) we shall denote by $p_V(y)$. It is well known that the number of linear independent equations is $m + n - 1$ and that the following holds

$$V \neq \emptyset \implies \sum_{i=1}^m a_i = \sum_{j=1}^n b_j. \quad (5)$$

At first, we shall find $p_{V_1}(y)$ and $p_{V_2}(y)$ explicitly. The solution of (P_{V_1}) is explicitly written as (Lawsdon, Hanson [74])

$$p_{V_1}(y) = y - A_1^T (A_1 A_1^T)^{-1} (A_1 y - a) \quad (6)$$

and if we introduce notation

$$\mu^* = (A_1 A_1^T)^{-1} (A_1 y - a)$$

we obtain $y - P_{V_1}(y) = A_1^T \mu^*$; μ^* is a vector of Lagrangian multipliers.

We denote by

$$\begin{aligned} A_1^+ &= A_1^T (A_1 A_1^T)^{-1} \\ A_2^+ &= A_2^T (A_2 A_2^T)^{-1} \end{aligned} \quad (7)$$

the Moore-Penrose generalized inverse A and by

$$\begin{aligned} Q_1 &= I - A_1^+ A_1 = I - P_1 \\ Q_2 &= I - A_2^+ A_2 = I - P_2 \end{aligned} \quad (8)$$

the projector onto the manifold defined by (2) and (3).

PROPOSITION 1.

$$\begin{aligned} A_1^+ &= (1/n)A_1^T \\ A_2^+ &= (1/m)A_2^T. \end{aligned}$$

Proof. It is easy to check that the following holds:

$$\begin{aligned} A_1 A_1^T &= nI, \\ A_2 A_2^T &= mI \end{aligned}$$

so we get from (7), $A_1^+ = (1/n)A_1^T$ and $A_2^+ = (1/m)A_2^T$.

Using Proposition 1, we get

$$\begin{aligned} P_{V_1}(y) &= y - (1/n)A_1^T(A_1 y - a) \\ P_{V_2}(y) &= y - (1/m)A_2^T(A_2 y - b) \end{aligned}$$

i.e.

$$[p_{V_1}(y)]_{ij} = y_{ij} - (1/n)(\sum y_{iq} - a_i) = y_{ij} - (u_i - a_i)/n \quad (9)$$

$$[p_{V_2}(y)]_{ij} = y_{ij} - (1/m)(\sum y_{pj} - b_j) = y_{ij} - (v_j - b_j)/m, \quad (10)$$

where $u_i = \sum_q y_{iq}$, $v_j = \sum_p y_{pj}$.

PROPOSITION 2. If we denote by $1_{p \times q}$ the matrix with p rows and q columns and with all elements equal to 1, then the following holds:

- 1° $A_2 A_1^T = 1_{n \times m}$
- 2° $A_1 A_2^T = 1_{m \times n}$
- 3° $(A_1 A_2^T) A_2 = 1_{m \times (mn)}$
- 4° $(A_2 A_1^T) A_1 = 1_{n \times (mn)}$
- 5° $A_2^T (A_2 A_1^T) = 1_{(mn) \times m}$
- 6° $A_1^T (A_1 A_2^T) = 1_{(mn) \times n}$
- 7° $A_2^T (A_2 A_1^T) A_1 = A_1^T (A_1 A_2^T) A_2 = 1_{(mn) \times (mn)}$.

Proof. Equalities 1° and 2° are obvious. Other equalities can be obtained from 1° and 2°.

PROPOSITION 3. $Q_2 Q_1 = Q_1 Q_2$.

Proof. After dividing 7° by mn , we get

$$(1/m)A_2^T (A_2 (1/n)A_1^T) A_1 = ((1/n)A_1^T) A_1 ((1/m)A_2^T) A_2.$$

From Proposition 1 and (8), we obtain

$$(A_2^+ A_2)(A_1^+ A_1) = (A_1^+ A_1)(A_2^+ A_2),$$

i.e.,

$$P_2 P_1 = P_1 P_2.$$

Now, we have ($Q_1 = I - P_1$ and $Q_2 = I - P_2$)

$$\begin{aligned} Q_2 Q_1 &= (I - P_2)(I - P_1) = I - P_1 - P_2 + P_2 P_1 = I - P_2 - P_1 + P_1 P_2 \\ &= (I - P_1)(I - P_2) - Q_1 Q_2. \end{aligned}$$

PROPOSITION 4. *If $V \neq 0$, then the following holds*

$$A_2^T A_2 A_1^T a = A_1^T A_1 A_2^T b. \quad (11)$$

Proof. Substituting 5° and 6° (from Proposition 2) in (11), we get

$$(1_{(mn)m})a = (1_{(mn)n})b,$$

which is equivalent with condition (5).

PROPOSITION 5. *For every $y \in R^{mn}$ the following holds*

$$p_{V_1}(p_{V_2}(y)) = p_{V_2}(p_{V_1}(y)).$$

Proof. From (6), (7) and Proposition 1, we obtain

$$\begin{aligned} p_{V_1}(p_{V_2}(y)) &= Q_1(p_{V_2}(y)) + A_1^+ a = Q_1(Q_2 y + A_2^+ b) A_1^+ a \\ &= Q_1 Q_2 y + (I - P_1) A_2^+ b + A_1^+ a = Q_1 Q_2 y + A_2^+ b + A_1^+ a - (1/mn) A_1^T A_1 A_2^T b. \end{aligned} \quad (12)$$

In the same way we have

$$p_{V_2}(p_{V_1}(y)) = Q_2 Q_1 y + A_2^+ b + A_1^+ a - (1/mn) A_2^T A_2 A_1^T a. \quad (13)$$

Because of Propositions 3 and 4, right-hand sides of (12) and (13) are equal, so Proposition 5 is proved.

Now, the projection point onto a manifold V could be achieved by consecutive projections onto V_1 and V_2 using simple formulas (9) and (10).

ALGORITHM

Step 1. There are given $y \in R^{mn}$, $a \in R^m$, $b \in R^n$

Step 2. Find $p_{V_1}(y)$ as $u_i = \sum_q y_{iq}$,

$$z_{ij} = y_{ij} - (1/n)(u_i - a_i), \quad i = 1, \dots, m; j = 1, \dots, n.$$

Step 3. Find $p_V(y)$ as $v_j = \sum_p z_{pj}$

$$x_{ij} = z_{ij} - (1/m)(v_j - b_j), \quad i = 1, \dots, n; j = 1, \dots, n.$$

3. CONCLUSION

Projection of a point onto a manifold needs $O(n^3)$ computations in general (in our case $O((m+n-1)^3)$). In this paper it is shown that this number is $O(m+n)$, for transportation manifold. For that purpose, the manifold is divided into two affine sets that have a very simple structure.

This result could be applied in various ways. For instance, in solving a linear transportation problem by "modified Karmarkar strategy" (Cavalier, Soyster [1985]), most time consuming activity is in finding projection of Dc onto the null space of AD ($D = \text{diag}(x_{11}, \dots, x_{mn})$, x — feasible point). Unfortunately, half-spaces $V'_1 = \{(A_1 D)_x = 0\}$ and $V'_2 = \{(A_2 D)_x = 0\}$ are not orthogonal, but an idea of consecutive projections onto these sets in order to find projections is still valid.

REFERENCES

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