

## ON THE DECOMPOSITION OF THE LINE (TOTAL) GRAPHS WITH RESPECT TO SOME BINARY OPERATIONS

*Slobodan K. Simić*

(Received March 8, 1978)

**Abstract:** It is proved that line and total graphs (and their complements too) are nearly always prime, i.e. excluding some exceptions, they could not be decomposed with respect to some binary operations (sum, product and strong product). In other words, the corresponding graph equations have not too many solutions.

### 0. Introduction

In this paper we shall consider only finite, undirected graphs without loops or multiple edges. Throughout the paper  $\bar{G}$ ,  $L(G)$ ,  $T(G)$ , as customary, will denote the complement, the line graph and the total graph of  $G$ . Further,  $\cup$  (in some papers  $+$ ) and  $\nabla$  (in [1] it is denoted by  $+$ ) will stand for union and join of graphs. We shall also prefer to use  $+$  for the cartesian product,  $\times$  for the conjunction and  $*$  for the strong product of some graphs. Next, we shall assume that two graphs  $G$  and  $H$  are equal if they are isomorphic

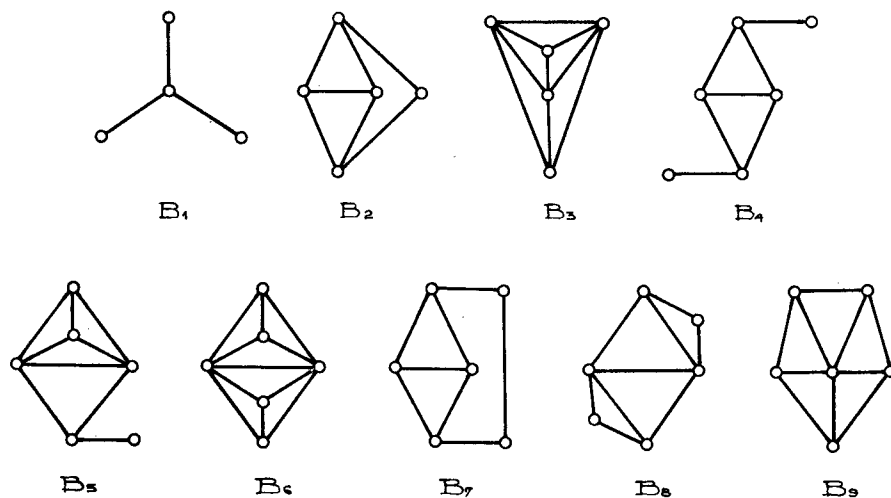


Fig 1.

and this will be denoted by  $G=H$ . If  $G$  is an induced subgraph of  $H$ , we shall write  $G\subseteq H$ . For the rest of notations, see [1]. Especially for the definition of graph equations, see [2] or [3].

For our further purposes we quote some facts about line and total graphs.

**Theorem 0.1.**  $G$  is a line graph if and only if none of the nine graphs from Fig. 1 is an induced subgraph of  $G$ .

**Theorem 0.2.** If  $G$  and  $H$  are connected graphs such that  $L(G)=L(H)$  then  $G=H$  unless one of them is  $K_3$  and the other is  $K_{1,3}$ .

These are well known theorems of L. W. Beineke and H. Whitney, for line graphs. There are also analogous results for total graphs (one can find them in [4], [5], [6]). Here we shall only give some lemmas which can be easily verified. Primarily, we need a concept of  $\nu$ -vertex and  $e$ -vertex of a total graph. If  $H=T(G)$  then some vertex of  $H$  is the  $\nu$ -vertex ( $e$ -vertex) if it originates from some vertex (edge) of  $G$ . Note, according to [4] if  $G$  has no cycles or complete graphs as components, then  $\nu$ -vertices or  $e$ -vertices are uniquely determined in  $H$ . Now, we give the following six lemmas, first three of which, can also be found in [7].

**Lemma 1.** In  $T(G)$  any  $e$ -vertex is adjacent to precisely two  $\nu$ -vertices.

**Lemma 2.** If in  $T(G)$  the  $e$ -vertex  $e$  is adjacent to  $\nu$ -vertices  $u$  and  $v$ , then  $u$  and  $v$  are adjacent.

**Lemma 3.** If in  $T(G)$  the  $\nu$ -vertex  $\nu$  is adjacent to  $e$ -vertices  $e$  and  $f$ , then  $e$  and  $f$  are adjacent.

**Lemma 4.** If  $K_{1,3}\subseteq T(G)$ , then the „central“ vertex of  $K_{1,3}$  is the  $\nu$ -vertex.

**Lemma 5.** If in  $T(G)$  two  $\nu$ -vertices are adjacent, then there is just one  $e$ -vertex adjacent to both of them.

**Lemma 6.** If in  $T(G)$  two  $e$ -vertices are adjacent, then there is just one  $\nu$ -vertex adjacent to both of them.

In Section *A* we shall treat the problems of the decomposition of the line graphs (and their complements too) according to binary operations referred as sum  $+$ , product  $\times$  and strong product  $*$ , while in Section *B* the analogous problems will be considered for total graphs. Some interesting remarks on total graphs are given in the Appendix.

## 1. Section A

In [8] or [9] one can find the following result, which is here restated in the form of graph equation.

**Theorem A1.** Graph equation  $L(G)=G_1+G_2$ , for  $G$  being connected and  $G_1, G_2$  being nontrivial, has the following solutions:  $(G, G_1, G_2) = (K_{m,n}, K_m, K_n)$ .

Of course, the restrictions given for graphs  $G, G_1, G_2$  do not, in essence, change the generality of the above result. In further text we shall use to make some similar restrictions in order to avoid any of uninteresting cases.

**Theorem A2.** *Graph equation  $\overline{L(G)} = G_1 + G_2$ , for  $G$  having no isolated vertices and for  $G_1, G_2$  being nontrivial, has the following solutions:*

$$(G, G_i, G_j) = (C_6, K_2, K_3), (K_{3,3}, K_3, K_3), (K_{1,mn}, mK_1, nK_1), (C_4, 2K_1, K_2),$$

$$(K_4, 3K_1, K_2), (2K_{1,2}, K_2, K_2), (4K_{1,2}, K_2, C_4), (K_{2,4} - 2K_2, K_2, K_{1,2})^*,$$

where  $i \neq j, i, j = 1, 2$ ,

**Proof.** Suppose first  $K_3 \subseteq G_i$ . But then  $G_j$  must be complete since otherwise  $K_1 \cup K_3 \subseteq 2K_1 + K_3 \subseteq G_1 + G_2$ , i.e.  $B_1 \subseteq L(G)$  (see Fig. 1). In the same way, it follows that  $G_j$  cannot have more than 3 vertices and that  $G_i$  cannot be a supergraph of  $K_3$  (otherwise  $K_1 \cup K_3$  appears again in  $G_1 + G_2$ ). Hence,  $(G_i, G_j) = (K_3, K_2), (K_3, K_3)$  and we get  $G$  easily.

Now let  $K_3 \not\subseteq G_1, G_2$ . Consequently  $K_3 \not\subseteq G_1 + G_2$ . Graphs  $H = \overline{L(G)}$  (for some  $G$ ) having no triangles were described in [10]. For our purpose the following observations are sufficient.

(a)  $K_{1,2}$  is not an induced subgraph for both  $G_1$  and  $G_2$ . Namely, if  $K_{1,2} \subseteq G_1, G_2$ , then  $K_2 \cup K_{1,2} \subseteq K_{1,2} + K_{1,2} \subseteq G_1 + G_2$ , i.e.  $B_2 \subseteq L(G)$  (see Fig. 1).

(b) If  $G_i$  has at least one edge, then  $P_4$  or  $K_{1,3}$  are not induced subgraphs of  $G_j$ . For  $P_4 \subseteq G_j$ , since  $K_2 \cup K_{1,2} \subseteq K_2 + P_4 \subseteq G_1 + G_2$ , we get the same contradiction as with (a). For  $K_{1,3} \subseteq G_j$ , then  $K_2 + K_{1,3}$  contains as an induced subgraph a graph obtained from a cycle of length 4 by adding two pendant edges to the same vertex of the cycle. Since the complement of the last graph is  $B_5$  (see Fig. 1) we arrive to a contradiction.

(c) If  $K_1 \cup K_2$  is an induced subgraph of  $G_i$ , then  $G_j = K_1$ . Namely, if  $G_j$  has at least two vertices, then  $K_2 \cup K_{1,2}$  or  $K_2 \cup 3K_1$  are induced subgraphs of  $G_1 + G_2$ . Since the complements of the above two graphs are  $B_2$  and  $B_3$  (see Fig. 1) we arrive to a contradiction.

Now we shall summarize the above observations. Clearly,  $G_1$  and  $G_2$  can be both without edges and then we find  $G$  easily. If only  $G_i$  is without edges then  $\overline{L(G)} = p(G_i)G_j$  ( $p(G_i) \geq 2$ ) and again if we want to avoid  $B_2 \subseteq L(G)$  or  $B_3 \subseteq L(G)$  it follows  $p(G_i) \leq 3$  and  $G_j = K_2$ . Thus,  $(G_i, G_j) = (2K_1, K_2), (3K_1, K_2)$  and for both cases  $G$  can be easily found. So, let both  $G_1, G_2$  have edges. Due to (c)  $G_1, G_2$  are both connected and since they are as well triangle-free according to (a) say  $G_i = K_2$  and from (b)  $G_j$  can be equal only to  $K_2, K_{1,2}, C_4$ . By direct checking we get that in all three possibilities  $G$  exists.  $\square$

\*  $K_{m,n} - pK_2$  is obtained from  $K_{m,n}$  by removing  $p$  independent edges.

**Theorem A 3.** *Graph equation  $L(G) = G_1 \times G_2$  for  $G$  having no isolated vertices and for  $G_1, G_2$  being connected and nontrivial, has the following solutions:*

$$(G, G_i, G_j) = (2P_{n+2}, P_{n+1}, P_2), (2C_{2n+2}, C_{2n+2}, K_2), (C_{4n+2}, C_{2n+1}, K_2), \\ (K_{3,3}, K_3, K_3), \text{ where } i \neq j, i, j = 1, 2.$$

**Proof.** Suppose that maximal vertex degree of  $G_1$  or  $G_2$  is greater than 2. Then, since  $G_1$  and  $G_2$  both have edges, it is easy too see that  $K_{1,3} \subseteq G_1 \times G_2$ , i.e.,  $K_{1,3} \subseteq L(G)$ . The last contradiction gives that  $G_1$  and  $G_2$ , due to connectedness, are paths or cycles. For  $K_{1,2} \subseteq G_i$  it follows that  $G_j = K_2$ , since otherwise (if  $G_j$  contains  $K_{1,2}$  or  $K_3$ )  $K_{1,3}$  appears again in  $G_1 \times G_2$ . Collecting the above conclusions we arrive at the proof of the theorem.  $\square$

The following two theorems are simple for proving and we only quote them.

**Theorem A 4.** *Graph equation  $\overline{L(G)} = G_1 \times G_2$ , for  $G$  having no isolated vertices and for  $G_1, G_2$  being nontrivial and not totally disconnected, has the following solutions:*

$$(G, G_1, G_2) = (K_{m,n}, K_m, K_n).$$

**Theorem A 5.** *Graph equation  $L(G) = G_1 * G_2$ , for  $G$  being connected and for  $G_1, G_2$  being nontrivial, has the following solutions:*

$$(G, G_1, G_2) = (K_{1,mn}, K_m, K_n).$$

**Theorem A 6.** *Graph equation  $\overline{L(G)} = G_1 * G_2$ , for  $G$  having no isolated vertices and for  $G_1, G_2$  being nontrivial, has the following solutions:*

$$(G, G_i, G_j) = (K_{1,mn}, mK_1, nK_1), (klK_2, K_k, K_l), (2n_1K_2 \cup n_2C_4 \cup n_3K_4, K_2, \\ \overline{n_1K_1 \cup n_2K_2 \cup n_3K_3}), \text{ where } i \neq j, i, j = 1, 2.$$

**Proof.** If  $G_1$  or  $G_2$  have no edges the same situation as with Theorem A 2 appears. So let both  $G_1$  and  $G_2$  have edges. If  $G_1$  and  $G_2$  are both complete we immediately get the corresponding graphs  $G$ . Thus, take that at least one of  $G_1$  and  $G_2$  is incomplete. Now, by observing that  $H_1 \subseteq G_1$  and  $H_2 \subseteq G_2$  implies  $H_1 * H_2 \subseteq G_1 * G_2$  we easily get that  $G_1$  and  $G_2$  have not  $K_2 \cup K_1$  as induced subgraphs, since otherwise  $K_1 \cup K_3 \subseteq (K_2 \cup K_1) * K_2$ , i.e.  $K_{1,3} \subseteq L(G)$  (note  $G_1$  and  $G_2$  both have edges). Hence, if  $G_j$  is incomplete it must contain  $K_{1,2}$  as an induced subgraph. But, then  $G_j$  cannot contain  $K_{1,2}$  or  $K_3$  as induced subgraphs since  $K_1 \cup K_3$  appears in  $K_{1,2} * K_{1,2}$  and  $K_3 * K_{1,2}$  as an induced subgraph. Thus  $G_i = K_2$ . For  $4K_1 \subseteq G_j$  we have  $K_2 \cup 3K_1 \subseteq G_1 * G_2$ , i.e.  $K_5 - x \subseteq L(G)$  (contradiction due to Theorem 0.1). Now we have  $G_j = n_1K_1 \cup n_2K_2 \cup n_3K_3$ . Since for each value of  $n_1, n_2, n_3$   $G$  exists we have found all the solutions.  $\square$

## 2. Section B

Now we shall prove the analogous theorems for totalgraphs.

**Theorem B 1.** *Graph equation  $T(G) = G_1 + G_2$ , for  $G_1, G_2$  being connected and nontrivial, has no solutions.*

**Proof.** Since both  $G_1, G_2$  have edges the same holds for  $G$ . But then  $T(G)$  contains a triangle and hence  $G_1$  or  $G_2$  must contain a triangle. Since  $G_1$  and  $G_2$  both have edges to each triangle in  $G_1 + G_2$  there corresponds another triangle such that they together induce a prism  $K_3 + K_2$ . Let us show that in  $T(G)$  there is a triangle without just mentioned property. Observe the triangle  $euv$  in  $T(G)$  such that  $e$  is an  $e$ -vertex and  $u, v$  are  $v$ -vertices. Suppose that  $x$  is a vertex of another triangle (which corresponds to the first one) and take that  $x$  is adjacent to  $e$  and not to  $u$  and  $v$ . By Lemma 1,  $x$  must be an  $e$ -vertex. Now, using Lemma 6 we easily get that  $x$  must be adjacent to  $u$  or  $v$ .  $\square$

**Theorem B 2.** *Graph equation  $\overline{T(G)} = G_1 + G_2$ , for  $G_1, G_2$  being non-trivial, has the following solutions:*

$$(G, G_i, G_j) = (K_3, 3K_1, K_2), \text{ where } i \neq j, i, j = 1, 2.$$

**Proof.** According to Theorem C 1 (see the Appendix)  $\overline{T(G)}$  is disconnected only if  $G = K_3$  or  $K_{1,n}$ . For  $G = K_3$  we easily get  $G_1$  and  $G_2$ , while for  $G = K_{1,n}$  there are no nontrivial solutions. Hence, we have that  $\overline{T(G)}$  is connected, what implies the same for  $G_1$  and  $G_2$ . If  $G_1$  and  $G_2$  are triangle-free the same holds for  $G_1 + G_2$ . In that case due to Theorem C 2 (see the Appendix)  $G_1 + G_2$  could be one of the following nine graphs:  $K_1, 3K_1, K_2, 3K_2, K_{1,3}, K_{3,3}, P_4 \cup K_1$ , Möbius ladder on eight vertices and Petersen's graph. By direct checking we confirm ourselves that in the above cases there are no additional solutions. So let  $G_i$  contain a triangle. Now, if we regard the vertices of  $G_1 + G_2$  as an ordered pairs then call  $V(k, x)$  ( $k = 1, 2$ ) the set of all vertices of  $G_1 + G_2$  having  $x$  for the  $k$ -th component. If  $G_j$  is incomplete, take  $V(j, x)$  and  $V(j, y)$  in such a way that  $x$  and  $y$  are not adjacent in  $G_j$ . In this case it is easy to see, by using Lemma 4, that all vertices in  $V(j, x) \cup V(j, y)$  regarded also as vertices of  $T(G)$  are  $v$ -vertices. Next, let  $u \in V(j, y)$  and take some  $e$ -vertex  $e$  (it must exist in  $T(G)$ ) adjacent to both  $u$  and  $v$ . By Lemma 1  $e$  is adjacent only to  $u$  and  $v$  among the vertices from  $V(j, x) \cup V(j, y)$ . But now in  $G_1 + G_2$  we then have that  $e \in V(j, z)$  ( $z \neq x, y$ ) is adjacent to all vertices in  $V(j, x)$  except  $u$ . Of course, this is a contradiction since  $|V(j, x)| = |V(G_j)| \geq 3$ . Therefore we have that  $K_3 \subseteq G_i$  implies  $G_j$  is complete, and accordingly, two cases are characteristic.

Case a:  $G_j$  has more than two vertices. Now  $K_3 \subseteq G_j$  and  $G_i$  is also complete. Then we have to find  $G$  such that  $\overline{T(G)} = K_m + K_n$  holds. Since  $K_m + K_n = \overline{K_m \times K_n}$  we have  $T(G) = K_m + K_n$ . Due to Theorem B 3 it follows that  $G$  does not exist.

Case b:  $G_j$  has just two vertices, i.e.  $G_j = K_2$ . Now  $G_i \neq K_3$  (see Case a) and since  $K_3 \subseteq G_i$  all vertices in  $V(j, x)$  and  $V(j, y)$  ( $x \neq y$ ) except perhaps three ones from each of the above sets are  $v$ -vertices (if regarded as vertices of  $T(G)$ ). The last assertion follows directly from Lemma 4. By Lemma 1, since  $G$  is not totally disconnected,  $p(G_i) - 3 \leq 2$ , i.e.  $p(G_i) \leq 5$ . Now it can be easily shown that even in this case there are no solutions.  $\square$

**Theorem B 3.** *Graph equation  $T(G) = G_1 \times G_2$ , for  $G_1, G_2$  being connected and nontrivial, has no solutions.*

**Proof.** Observing the maximal vertex degrees of  $T(G)$  and  $G_1 \times G_2$ , by using the relations  $\Delta(T(G)) = 2\Delta(G)$  and  $\Delta(G_1 \times G_2) = \Delta(G_1)\Delta(G_2)$ , we get

$$(a) \quad 2\Delta(G) = \Delta(G_1)\Delta(G_2).$$

Similarly, observing the number of vertices in the largest clique of the corresponding graphs we have  $c(T(G)) \geq \Delta(G) + 1$  and  $c(G_1 \times G_2) \leq \min(c(G_1), c(G_2)) \leq \min(\Delta(G_1), \Delta(G_2)) + 1$ . Therefore, it follows

$$(b) \quad 2\Delta(G) \leq \Delta(G_1) + \Delta(G_2).$$

Combining (a) and (b) we get

$$(c) \quad (\Delta(G_1) - 1)(\Delta(G_2) - 1) \leq 1.$$

For  $\Delta(G_i) = 1$  ( $i = 1, 2$ ), we have  $G_i = K_2$  (since  $G_i$  connected) and now  $G_1 \times G_2$  is bichromatic while  $T(G)$  is never bichromatic. So, by using (c) we get  $\Delta(G_1) = \Delta(G_2) = 2$ . Since  $G_1$  and  $G_2$  are connected they can be only paths or cycles and moreover because they cannot be bichromatic they are cycles of odd length, i.e.  $G_1 = C_{2m+1}$ ,  $G_2 = C_{2n+1}$ . Since  $C_{2m+1} \times C_{2n+1} = C_{2m+1} + C_{2n+1}$  according to Theorem B 1 there are no solutions.  $\square$

**Theorem B 4.** Graph equation  $\overline{T(G)} = G_1 \times G_2$ , for  $G_1, G_2$  being non-trivial, has no solutions.

**Proof.** Since  $\overline{T(G)}$  is disconnected only for  $G = K_3$  or  $K_{1,n}$  (see Theorem C 1) it immediately follows that  $G_1$  or  $G_2$  cannot be disconnected. As in Theorem B 2, if we regard the vertices of  $G_1 \times G_2$  as ordered pairs then call  $V(k, x)$  ( $k = 1, 2$ ) the set of all vertices of  $G_1 \times G_2$  having  $x$  for the  $k$ -th component. Now in  $G_1 \times G_2$  all vertices in  $V(i, x)$  ( $i = 1, 2$ ) are mutually non-adjacent. In  $\overline{G_1 \times G_2}$  the corresponding vertices are all mutually adjacent. Hence,  $\overline{G_1 \times G_2}$  can be partitioned into  $p(G_i)$  complete subgraphs with  $p(G_j)$  ( $j \neq i, j = 1, 2$ ) vertices. If  $p(G_j) \geq 4$  the vertices of the mentioned complete subgraphs regarded also as vertices of  $T(G)$  can be according to Lemma 5, only of the following types:

(a) all vertices in one complete subgraph are  $v$ -vertices;  
 (b) all vertices in one complete subgraph except one which is a  $v$ -vertex are  $e$ -vertices;

(c) all vertices in one complete subgraph are  $e$ -vertices.

Assume that in  $\overline{G_1 \times G_2}$  there are no complete subgraph of the type (a). Then we immediately get a contradiction since in that case  $G$  would have more edges than a complete graph with the same number of vertices. Now suppose that more than one of such complete subgraphs exist and observe just two having for their vertex sets  $V(i, x)$  and  $V(i, y)$  ( $x \neq y$ ). Next let an  $e$ -vertex  $e$  be adjacent to two (adjacent) vertices from  $V(i, x)$ . By Lemma 1 it cannot be adjacent to any vertex from  $V(i, y)$ . Hence, in  $G_1 \times G_2$  there exists a vertex being adjacent to all vertices from  $V(i, x)$  but not to all vertices from  $V(i, y)$  (note  $p(G_j) \geq 4$ ). Therefore, it follows that

only one complete subgraph of the type (a) may exist in  $\overline{G_1 \times G_2}$ . Denote it by  $C^1$  ( $V(i, x)$  is its vertex-set) and observe an  $e$ -vertex  $e$  adjacent to two vertices, say  $u, v$  from  $C^1$ . Of course,  $e$  belongs to some other complete subgraph, say  $C^2$  having  $V(i, y)$  as its vertex-set, and due to Lemma 1  $C^2$  is of the type (c). Hence, in  $T(G)$  due to Theorem 0.2 there is a star as a subgraph (not necessarily induced) having  $p(G_j)$  edges so that all its vertices are  $v$ -vertices and  $u$  or  $v$  is its "central" vertex. Now since  $C^1$  contains  $p(G_j)$  vertices at least one of the vertices of the mentioned star does not belong to  $C^1$ . Denote one such vertex by  $w$  and take that it belongs to some complete subgraph  $C^3$  (let the vertices of  $C^3$  be  $V(i, z)$ ,  $z \neq x, y$ ). As follows from above,  $C^3$  must be of the type (b). Now take a vertex  $f$  from  $C^2$  which is adjacent to  $w$ . Using Lemma 2 it follows that  $f$  is adjacent to all  $e$ -vertices from  $C^3$ . But now in  $G_1 \times G_2$   $f$  is not adjacent to all vertices in  $V(i, z)$  while  $e$  is adjacent to  $w \in V(i, z)$  and also we have  $e, f \in V(i, y)$ . The last is possible only if  $G_1$  or  $G_2$  have isolated vertices which contradicts the fact that  $G_1$  and  $G_2$  are connected. If  $p(G_i) \leq 3$  ( $i = 1, 2$ ), by direct checking, we get that there are no solutions.  $\square$

**Theorem B 5.** *Graph equation  $T(G) = G_1 * G_2$ , for  $G_1, G_2$  being connected and nontrivial, has no solutions.*

**Proof.** The case when  $G$  is without edge is excluded by suppositions of the theorem. Hence,  $T(G)$  contains a triangle and also  $G_1 * G_2$ . Now we shall prove that to each triangle of  $G_1 * G_2$  there corresponds in the same graph a vertex adjacent to all vertices of the observed triangle. For that purpose denote by  $(u_1, v_1), (u_2, v_2), (u_3, v_3)$  the vertices of some triangle in  $G_1 * G_2$  where  $u_1, u_2, u_3$  ( $v_1, v_2, v_3$ ) are the vertices of  $G_1$  ( $G_2$ ). Then one of the following possibilities must hold:

- (a)  $u_1, u_2, u_3$  and  $v_1, v_2, v_3$  form triangles in  $G_1$  and  $G_2$ , respectively;
- (b)  $u_1, u_2, u_3$  (or  $v_1, v_2, v_3$ ) form a triangle in  $G_1$  ( $G_2$ ) and just two vertices among  $v_1, v_2, v_3$  (or  $u_1, u_2, u_3$ ) say  $v_1, v_2$ , (or  $u_1, u_2$ ) represent the same vertex of  $G_2$  ( $G_1$ ) while the third one  $v_3$  (or  $u_3$ ) is adjacent to them;
- (c)  $u_1, u_2, u_3$  (or  $v_1, v_2, v_3$ ) form a triangle in  $G_1$  ( $G_2$ ) and all  $v_1, v_2, v_3$  (or  $u_1, u_2, u_3$ ) represent the same vertex in  $G_2$  ( $G_1$ );
- (d) two vertices among  $u_1, u_2, u_3$ , say  $u_1, u_2$  and also two vertices among  $v_1, v_2, v_3$  say  $v_2, v_3$  represent the same vertices in  $G_1$  and  $G_2$  respectively, while the third ones  $u_3$  and  $v_1$  are adjacent to  $u_1, u_2$  and  $v_2, v_3$ .

Now we shall find  $(u, v)$  adjacent to  $(u_i, v_i)$  ( $i = 1, 2, 3$ ).

Case a: It is sufficient to take  $u = u_i, v = v_j$   $i \neq j, i, j \in \{1, 2, 3\}$ .

Case b: It is sufficient to take  $u = u_3$  (or  $v = v_3$ ) and  $v = v_1 = v_2$  (or  $u = u_1 = u_2$ ).

Case c: It is sufficient to take  $u = u_i$  (or  $v = v_i$ )  $i \in \{1, 2, 3\}$  and  $v$  (or  $u$ ) is an arbitrary vertex of  $G_2$  ( $G_1$ ) adjacent to  $v_i$  (or  $u_i$ ).

Case d: Now take  $u = u_3$  and  $v = v_1$ .

To prove the theorem observe a triangle in  $T(G)$  having one  $e$ -vertex and two  $v$ -vertices. Such a triangle has not the above property since otherwise we get contradictions to Lemmas 1 and 5.  $\square$

**Theorem B 6.** Graph equation  $\overline{T(G)} = G_1 * G_2$ , for  $G_1, G_2$  being non-trivial, has the following solutions:

$(G, G_i, G_j) = (mnk_1, K_m, K_n), (2kK_1 \cup K_3, K_2, kK_1 \cup K_3)$ , where  $i \neq j, i, j = 1, 2$ .

**Proof.** It is easy to see that both  $G_1$  and  $G_2$  can be complete graphs in which case  $G$  is totally disconnected. Hence, suppose first that only  $G_i$  is complete and also let it be different from  $K_2$  while  $G_j$  is incomplete. If the vertices of  $G_1 * G_2$  are regarded as ordered pairs let, as earlier,  $V(k, x)$  ( $k = 1, 2$ ) denote all vertices having  $x$  for the  $k$ -th component. Now in  $G_1 * G_2$  vertices  $V(j, x)$  for all  $x$  are mutually nonadjacent and if  $u \in V(j, x)$  then  $u$  is adjacent or nonadjacent to all vertices belonging to  $V(j, y)$  for some  $y \neq x$ . So, each vertex in  $G_1 * G_2$  is isolated or adjacent to three mutually nonadjacent vertices (since  $p(G_i) \geq 3$ ). If the vertices of  $G_1 * G_2$  are now regarded also as the vertices of  $T(G)$  by using Lemma 4 it follows that they are all  $v$ -vertices and this is obviously a contradiction. Thus  $G_1$  and  $G_2$  are incomplete or one of them is equal to  $K_2$ .

For the beginning let both  $G_1$  and  $G_2$  be different from  $K_2$ . If  $G_i$  contains a triangle then take  $V(j, x)$  and  $V(j, y)$  so that  $x$  and  $y$  are nonadjacent. In  $G_1 * G_2$  each vertex from  $V(j, x)$  is adjacent to each vertex from  $V(j, y)$  and since  $G_i$  contains a triangle by using Lemma 4 it follows that all vertices in  $V(j, x) \cup V(j, y)$  are  $v$ -vertices. Since  $G_i$  is incomplete we can take in  $G_1 * G_2$  two adjacent vertices  $u, v \in V(j, x)$  and an  $e$ -vertex  $e \in V(j, z)$  ( $z \neq x, y$ ) which is adjacent to both  $u$  and  $v$ . By Lemma 1  $e$  is adjacent only to  $u$  and  $v$  among vertices belonging to  $V(j, x) \cup V(j, y)$ . Now in  $G_1 * G_2$   $e$  is adjacent to all except two vertices from  $V(j, x)$  and also it is adjacent to all vertices from  $V(j, y)$ . This is impossible by the definition of the operation  $*$  and thus  $G_1$  and  $G_2$  are both triangle-free. Assume now that  $K_1 \cup K_2 \subseteq G_i$ . Since  $G_j$  is not complete or totally disconnected (the last follows from Theorem C 1) in  $G_j$  there exists a vertex being adjacent and nonadjacent to some other vertices of  $G_j$ . Hence, it is possible to find in  $G_1 * G_2$  three groups of vertices, say  $V(j, x), V(j, y), V(j, z)$  ( $x, y, z$  are mutually different), so that  $x$  and  $y$  are nonadjacent and  $x$  and  $z$  are adjacent. Next, let  $u_t = (u, t), v_t = (v, t), w_t = (w, t)$  ( $t = x, y, z$ ) and according to the above we can choose that  $u_t, v_t, w_t$  induce  $K_1 \cup K_2$  in  $G_1 * G_2$ . Now in  $G_1 * G_2$  we have the following induced subgraph, see Fig 2.

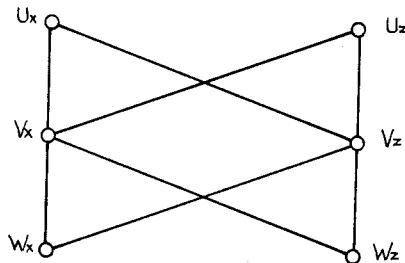


Fig 2.



Now if we regard vertices  $u_t, v_t, w_t$  ( $t = x, y, z$ ) as vertices of  $T(G)$  then by Lemma 4 (see Fig. 2)  $v_x$  and  $v_z$  must be  $v$ -vertices and by Lemma 2  $u_x, w_x, u_z, w_z$  must be also  $v$ -vertices. Hence, all vertices in  $V(j, y)$  (see Lemma 1) must be  $v$ -vertices and consequently all vertices in  $V(j, x)$  must be  $v$ -vertices. Choose an  $e$ -vertex  $e$  adjacent to  $u_x, v_x$ . By Lemma 1 it is nonadjacent to all other vertices from  $V(j, x) \cup V(j, y)$ . So in  $G_1 * G_2$   $e$  ( $e \notin V(j, x) \cup V(j, y)$ ) is adjacent to all vertices except two from  $V(j, x)$  and also it is adjacent to all vertices from  $V(j, y)$ . The last is, of course, a contradiction and thus  $K_1 \cup K_2$  is not an induced subgraph in  $G_1$  and  $G_2$ . Now since  $K_1 \cup K_2$  and  $K_3$  are both forbidden in  $G_1$  and  $G_2$  as induced subgraphs and also since neither of  $G_1$  and  $G_2$  is totally disconnected (see Theorem C1)  $G_1$  and  $G_2$  must be bicomplete graphs. But then it is a matter of routine, by using Lemmas 1 — 6, to verify that the corresponding graph  $G$  does not exist. Namely, if  $G_1 = K_{m,n}$  and  $G_2 = K_{k,l}$  from the above mentioned lemmas for  $\max(m, n, k, l) \geq 3$ , we easily get contradictions; for the rest of possibilities by direct checking we find that  $G$  does not exist.

At last it remains the case  $G_i = K_2$  and  $G_j$  is incomplete. If  $K_1 \cup K_2 \subseteq G_j$  then we have the same situation as earlier (see Fig. 2 and corresponding conclusions) i.e. now all vertices in  $\overline{G_1 * G_2}$ , if regarded as vertices of  $T(G)$ , are  $v$ -vertices, and this is obviously absurd. Hence,  $K_1 \cup K_2$  is forbidden in  $G_j$  as an induced subgraph and so  $\overline{G_j} = \bigcup_i m_i K_i$ . Now  $\overline{G_1 * G_2} = \bigcup_i m_i (iK_2)$  and since this graph must also be a total graph all  $m_i$  except  $m_1$  and  $m_3$  are equal to 0.  $\square$

### 3. Appendix

Here we will prove two theorems which are used in the above text and which also give some analogous results for complementary total graphs in comparison with the results obtained in [10], for complementary line graphs.

**Theorem C1.** *If  $H = \overline{T(G)}$  for some  $G$ , then  $H$  is disconnected if and only if  $G = K_3$  or  $K_{1,n}$ .*

**Proof.** In fact we have to solve a graph equation  $T(G) = G_1 \nabla G_2$ . Since  $G$  is not totally disconnected we have an  $e$ -vertex  $e \in V(G_i)$  ( $i = 1$  or  $2$ ). But now, since  $e$  is adjacent to all vertices from  $V(G_j)$  ( $j \neq i, j = 1, 2$ ), by using Lemma 1, we get that  $V(G_j)$  has at most two  $v$ -vertices.

Case a:  $G_j$  contains just two (adjacent)  $v$ -vertices. Due to Lemma 5 all vertices in  $G_i$  except  $e$  (if any exists) are  $v$ -vertices. If  $G_i$  contains only the vertex  $e$  we get  $G = K_2$ . Otherwise in  $G_j$  there must exist at least two  $e$ -vertices but due to Lemma 5  $G$  can contain only one  $v$ -vertex. Thus, we get  $G = K_3$ .

Case b:  $G_j$  contains only one  $v$ -vertex, say  $u$ . Then  $G$  contains a  $v$ -vertex  $v$  adjacent to  $e$ . If  $G_j$  contains at least one  $e$ -vertex then, due to Lemma

5, in  $G_i$  exist at least two  $v$ -vertices and we get the same situation as above. So let  $u$  be the only vertex in  $G_j$ . Suppose now that in  $G_i$  exists a  $v$ -vertex  $w$  adjacent to  $v$ . This is due to Lemma 1 impossible, since otherwise an  $e$ -vertex  $f$  would be adjacent to  $u$ ,  $v$  and  $w$ . If all  $v$ -vertices from  $G_i$  are mutually nonadjacent we get  $G = K_{1,n}$ .

Case c:  $G_j$  has no  $v$ -vertices. Now it is easy to see that except  $G = K_2$  there are no other possibilities.  $\square$

Corollary. Graph equation  $T(G) = G_1 \nabla G_2$  has only the following solutions:

$$(G, G_i, G_j) = (K_3, 2K_1, C_4), (K_{1,n}, K_1, K_n \circ K_1), \text{ where } i \neq j, i, j = 1, 2.$$

Here  $\circ$  denotes the graph operation corona, see [1].

Theorem C2. A graph  $G$  without triangles is a complement of a total graph if and only if it is equal to one of the following nine graphs:  $K_1$ ,  $3K_1$ ,  $K_2$ ,  $3K_2$ ,  $K_{1,3}$ ,  $K_{3,3}$ ,  $P_4 \cup K_1$ , Möbius ladder on eight vertices and Petersen's graph.

Proof. Observe that  $H_0 \subseteq H$  implies  $\overline{T(H_0)} \subseteq \overline{T(H)}$ . Now it is easy to see that  $3K_1 \subseteq \overline{T(H_0)}$  for each  $H_0 \in \{3K_1, K_3 \cup K_1, P_4, K_{1,3} + x, K_4 - x, K_5\}$ . Thus all graphs from the above set are forbidden as induced subgraphs of  $H$  if  $T(H)$  should be triangle-free. Now it is a matter of routine to find all graphs  $H$  (just nine) having the above property. Since their complementary total graphs are triangle-free, the theorem is proved.  $\square$

#### REFERENCES

- [1] F. Harary, *Graph theory*. Reading 1969.
- [2] D. M. Cvetković, I. B. Lacković, S. K. Simić, *Graph equations, graph inequalities and a fixed point theorem*, Publ. Inst. Math. (Beograd), 10 (34) (1976), 59–66.
- [3] D. M. Cvetković, S. K. Simić, *Graph equations, Beiträge zur Graphentheorie und deren Anwendungen, vorgetragen auf dem Internat. Koll. Oberhof (DDR)*, 10–16. April 1977, 40–56.
- [4] M. Behzad, H. Radjavi, *The total group of a graphs*, Proc. Amer. Math. Soc. 19 (1968), 158–163.
- [5] M. Behzad, H. Radjavi, *Structure of regular total graphs*, J. London Math. Soc. 44 (1969), 433–436.
- [6] M. Behzad H. *A characterization of total graphs*, Proc. Amer. Math. Soc. 26 (1970), 383–389.
- [7] F. Escalante, J. M. S. Simões-Pereira, *Just two total graphs are complementary*, Monatshefte für Math. 81 (1976), 5–13.
- [8] E. M. Palmer, *Prime line graphs*, Nanta Math. 6 (1973), No. 2, 75–76.
- [9] M. Doob, *A note on prime graphs*, Utilitas Math. 9 (1976), 297–299.
- [10] D. M. Cvetković, S. K. Simić, *Some remarks on the complement of a line graph*, Publ. Inst. Math. (Beograd), 17 (31) (1974), 37–44.