

## ON THE EMBEDDING OF $\Omega$ -ALGEBRAS IN GROUPOIDS

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**Summary.** It has been proved that any  $\Omega$ -algebra can be embedded both in a semigroup [1], [3], and in a so called entropic groupoid [2].

This paper gives a necessary and sufficient condition for embedding any  $\Omega$ -algebra in some groupoid  $\mathcal{G} = (G, *)$  satisfying<sup>1)</sup> the set of laws  $\Sigma(*)$ . The condition is:

*There is a term  $\xi(x, y; *)$  formed of two variables  $x, y$  and the operation symbol  $*$  such that the operation  $\circ$  defined by*

$$(E) \quad xy \circ \stackrel{\text{def}}{=} \xi(x, y; *)$$

*does not satisfy any algebraic law (except the law  $x = x$ ), while the operation  $*$  satisfies the laws in  $\Sigma(*)$ .*

Some new examples of groupoids which satisfy the condition (E) are presented:

1° Each  $\Omega$ -algebra can be embedded in a commutative groupoid.

2° Each  $\Omega$ -algebra can be embedded in a groupoid satisfying the law of the type

$$\Pi x_1 \cdots x_n = \Pi x_{p_1} \cdots x_{p_n},$$

where at both sides of the equality the arrangement of the operation symbols is the same, and the permutation  $\begin{pmatrix} 1 & 2 & \cdots & n \\ p_1 & p_2 & \cdots & p_n \end{pmatrix}$  has at least one fixed point.

For example, the laws of this type are

$$xy * z * = zy * x *, \quad xy * u * v w * * = w y * v * u x * *.$$

In the second part of the paper it is shown that the condition (E) can be extended so that the main theorems 1, 2 hold for the laws  $\Sigma(O)$ , i. e. for variety  $V_O(\Sigma)$ , where  $O$  is a set of some operation symbols. For example, any  $\Omega$ -algebra can be embedded in an entropic algebra, i. e. in an algebra satisfying:

$$(x_{11}x_{12} \cdots x_{1n}f \cdots x_{n1}x_{n2} \cdots x_{nn}f)f = (x_{11}x_{21} \cdots x_{n1}f \cdots x_{1n}x_{2n} \cdots x_{nn}f)f$$

<sup>1)</sup> In other words,  $\mathcal{G}$  belongs to the variety  $V_*(\Sigma)$ , [1].

1. The main result of the paper is the following

**Theorem 1.** *If  $Q$  is an arbitrary  $\Omega$ -algebra and  $\Sigma(*)$  a set of laws<sup>1)</sup> satisfying the condition (C), then there exists a groupoid  $\mathcal{G}=(G, *)$  satisfying the laws  $\Sigma(*)$ , with the properties:*

(i)  $Q$  is a subset of  $G$ ;

(ii) *If  $\omega \in \Omega$  is some  $n$ -ary operation, then there exists an  $\bar{\omega} \in G$  and a term  $t_\omega(x_1, \dots, x_n, \bar{\omega}; \circ)$  formed from variables  $x_1, \dots, x_n$ , the constant symbol  $\bar{\omega}$  and the operation symbol  $\circ$  such that:*

$$\omega = \bar{\omega}, \quad \text{if } \omega \in \Omega(\circ)$$

$$x_1 x_2 \cdots x_n \omega = t_\omega(x_1, \dots, x_n, \bar{\omega}; \circ), \quad \text{if } \omega \in \Omega(n), \quad n = 1, 2, \dots$$

where  $xy \circ \stackrel{\text{def}}{=} \xi(x, y; *)$ . The existence of  $\xi(x, y; *)$  is guaranteed by (C).

For the proof we need the following definitions:

I. Let  $G$  be a minimal set satisfying the conditions:

The set  $X \stackrel{\text{def}}{=} Q \cup \Omega$  is a subset of  $G$ ,

If  $u, v \in G$ , then  $uv * \in G$ .

Let  $\mathcal{G}=(G, *)$  be a groupoid determined by the set  $G$  and by the operation defined in the natural way, i. e.  $\mathcal{G}$  is the  $*$ -word algebra on  $X$ .

II. Let  $G_0$  be a minimal subset of  $G$  such that:

$X$  is a subset of  $G_0$ ,

If  $u, v \in G_0$ , then  $uv \circ \in G_0$ , where  $\circ$  is defined by the term  $\xi(x, y; *)$  (of condition (C)):

$$\text{Def } (\circ) \quad xy \circ \stackrel{\text{def}}{=} \xi(x, y; *)$$

III. For each  $\omega \in \Omega(n)$  ( $n = 0, 1, \dots$ ) we define an operation  $\omega$  on  $G$ :

$$\text{Def } (\Omega) \quad \omega = \omega, \quad \text{if } n = 0,$$

$$x_1 \dots x_n \omega = t_\omega(x_1, \dots, x_n, \omega; \circ), \quad \text{if } n = 1, 2, \dots,$$

where  $t_\omega(x_1, \dots, x_n, \omega; \circ)$  is a term formed from variables  $x_1, \dots, x_n$ , the constant  $\omega$  and the operation symbol  $\circ$ . These terms can be chosen arbitrarily, in particular, they can be chosen with the operation symbols grouped on the right [2]. The operations defined in that way form the set

$$\Omega \stackrel{\text{def}}{=} \{\omega \mid \omega \in \Omega\}.$$

IV. Finally, let  $G_\Omega$  be a minimal subset of  $G_0$  such that:

$Q \cup \Omega(0)$  is a subset of  $G_\Omega$ ,

If  $u_1, \dots, u_n \in G_\Omega$ , and  $\omega \in \Omega(n)$  ( $n = 1, 2, \dots$ ), then  $u_1 \dots u_n \omega \in G_\Omega$ . Hence, the elements of  $G_\Omega$  are those terms of  $G$  which can be represented by operations from  $\Omega$ .

<sup>1)</sup> We assume that the laws are consistent, i. e. that they do not imply  $x=y$ , where  $x, y$  are different variables.

V. Let  $\sim_Z, \sim_Q, \sim_\circ, \sim_\Omega$  be the *minimal congruences* generated by

$$\Sigma(*), \text{Tab } Q(\Omega), \text{Def}(\circ), \text{Def}(\Omega),$$

that is:

$$\begin{aligned} u \sim_Z v &\stackrel{\text{def}}{\Leftrightarrow} \Sigma(*) \vdash u = v, \\ u \sim_Q v &\stackrel{\text{def}}{\Leftrightarrow} \text{Tab } Q(\Omega) \vdash u = v, \quad (u, v \in G) \\ u \sim_\circ v &\stackrel{\text{def}}{\Leftrightarrow} \text{Def}(\circ) \vdash u = v \\ u \sim_\Omega v &\stackrel{\text{def}}{\Leftrightarrow} \text{Def}(\Omega) \vdash u = v. \end{aligned}$$

The symbol  $\vdash$  denotes the *logical deduction* [4], for example  $\Sigma(*) \vdash u = v$  means that the formula  $u = v$  can be derived from the laws  $\Sigma(*)$  and the equality axioms.  $\text{Tab } Q(\Omega)$  is the so called *positive diagram* [4] of the algebra  $Q$ , i. e. the set of all equalities of the form

$$a_1 \cdots a_n \omega = a \quad (a_1, \dots, a_n, a \in Q, \omega \in \Omega(n), n = 1, 2, \dots)$$

which hold in  $Q$ .  $\text{Tab } Q(\Omega)$  is the corresponding set of formulas obtained by exchanging each  $\omega \in \Omega$  by  $\omega \in \Omega$ .

VI. The relation  $\sim$  is the *minimal congruence* of the set  $G$  generated by

$$\sim_Z, \sim_Q, \sim_\circ, \sim_\Omega.$$

*Lemma.* Let  $\rho$  be one of the relations  $\sim_Z, \sim_Q, \sim_\circ, \sim_\Omega$  and let  $\sigma$  be one of the relations  $\sim_Z, \sim_\circ, \sim_\Omega$ . Then

$$(u \in G_\Omega \wedge u \rho v) \Rightarrow (\exists v' \in G_\Omega) v \sigma v'.$$

*Proof.* We distinguish four cases:  $u \sim_Z v, u \sim_Q v, u \sim_\circ v, u \sim_\Omega v$ .

Ad1. In this case the term  $v'$  is just  $u$ , since by condition (C), each term can be uniquely represented by  $\circ$  (if such a representation exists), and hence uniquely represented by operations from  $\Omega$  (if the representation exists).

Ad2. Let  $u = u(\cdots a_1 \cdots a_n \omega \cdots)$ ,  $a_1, \dots, a_n \in Q$ . If the subterm of  $u$ :  $a_1 \cdots a_n \omega$  is replaced by  $a$ , under the condition that (in  $Q$ )  $a_1 \cdots a_n \omega = a$ , the resulting term is again in  $G_\Omega$ . Similarly, in the case  $u(\cdots a \cdots)$ , i. e.  $a$  is a subterm of  $u$ , we conclude that  $u(\cdots a_1 \cdots a_n \omega \cdots) \in G_\Omega$ , if  $a_1 \cdots a_n \omega = a$  holds in  $Q$ . Since  $\sim_Q$  can be defined in a finite number of such replacements, we have

$$u \in G_\Omega \wedge u \sim_Q v \Rightarrow v \in G_\Omega.$$

In cases 3 and 4, it immediately follows that  $v'$  is uniquely determined and equal to  $u$ . Hence:

- (1) *In the cases  $u \sim_Z v, u \sim_\circ v, u \sim_\Omega v$ , if  $u \in D_\Omega$ , then the uniquely determined  $u'$  is just  $u$ .*

*Proof of Theorem 1.* We first prove:

- (2)  $\bar{x} = \bar{y} \Rightarrow x = y \quad (x, y \in Q),$

where  $\bar{x}, \bar{y}$  are the equivalence classes of  $x, y$  with respect to  $\sim$ .

Suppose that  $\bar{x} = \bar{y}$ , i. e.  $x \sim y$ . Then there exists a natural number  $k$  and elements

$$u_1, u_2, \dots, u_k; \quad u_1 = x, \quad u_k = y \quad (u_i \in G)$$

such that for each  $i = 1, \dots, k-1$ :

$$u_i \sim_Z u_{i+1} \text{ or } u_i \sim_Q u_{i+1} \text{ or } u_i \sim_\circ u_{i+1} \text{ or } u_i \sim_\Omega u_{i+1}.$$

Let **Int** be an interpretation, i. e. a mapping such that

- (i) If  $x \in Q$ , then  $\mathbf{Int}(x) \stackrel{\text{def}}{=} x$ ,
- (ii) If  $\omega \in \Omega(0)$ , then  $\mathbf{Int}(\omega) \stackrel{\text{def}}{=} \omega$ ,
- (iii) If  $\omega \in \Omega(n)$  ( $n = 1, 2, \dots$ ) and  $t_1, \dots, t_n \in D_\Omega$ , then

$$\mathbf{Int}(t_1 \cdots t_n \omega) \stackrel{\text{def}}{=} \mathbf{Int}(t_1) \cdots \mathbf{Int}(t_n) \omega$$

- (iv) If  $t \in D$  and there exists a  $t' \in D_\Omega$  such that

$$\Sigma(*), \text{ Def}(\circ), \text{ Def}(\Omega) \vdash t = t'$$

then  $\mathbf{Int}(t) \stackrel{\text{def}}{=} \mathbf{Int}(t')$ .

The mapping **Int** is well defined. The conditions (i), (ii), (iii) represent the usual definition of homomorphism. The soundness of part (iv) follows from (1).

The mapping **Int** carries the sequence  $u_1, \dots, u_k$  into the sequence

$$\mathbf{Int}(u_1), \mathbf{Int}(u_2), \dots, \mathbf{Int}(u_k)$$

of elements of  $\Omega$ -algebra  $Q$ :

From the definition of the mapping **Int** we have:

If one of the conditions  $u \sim_Z v$ ,  $u \sim_Q v$ ,  $u \sim_\circ v$ ,  $u \sim_\Omega v$  is satisfied, then the equality  $\mathbf{Int}(u) = \mathbf{Int}(v)$  holds in  $Q$ .

Hence:

$$x = \mathbf{Int}(u_1) = \mathbf{Int}(u_2) = \dots = \mathbf{Int}(u_k) = y$$

implying that the equality  $x = y$  holds in  $Q$ , which proves (2).

In order to complete the proof of the theorem, we introduce the quotient groupoid  $\bar{\mathcal{G}} = (\bar{G}, \bar{*})$ , where

$$\bar{G} \stackrel{\text{def}}{=} \{\bar{x} \mid x \in G\}, \quad \bar{x}\bar{y}\bar{*} \stackrel{\text{def}}{=} \overline{xy*}.$$

The groupoid  $\bar{\mathcal{G}}$  satisfies the same laws as  $\mathcal{G}$ .

Further, let

$$\bar{Q} \stackrel{\text{def}}{=} \{\bar{x} \mid x \in Q\}.$$

For each  $\omega \in \Omega(n)$  we define an operation  $\bar{\omega}$  in  $\bar{Q}$  in the following way

$$\bar{\omega} \stackrel{\text{def}}{=} \bar{\omega}, \quad \text{if } \omega \in \Omega(0)$$

$$\bar{x}_1 \cdots \bar{x}_n \bar{\omega} \stackrel{\text{def}}{=} \overline{x_1 \cdots x_n \omega}, \quad \text{if } \omega \in \Omega(n) \quad (n = 1, 2, \dots).$$

It is clear that  $\bar{Q}$  is an  $\Omega$ -algebra ( $\bar{\omega}$  corresponds to  $\omega$ ). From the first part of the proof it follows that  $Q$  and  $\bar{Q}$  are isomorphic algebras (an isomorphism being  $f: x \rightarrow \bar{x}$ ). This completes the proof of the theorem.

2. We give some examples of the laws  $\Sigma(*)$  satisfying the condition ( $\mathcal{C}$ ).

I. Already known examples of the laws  $\Sigma(*)$  are associative and entropic laws ([1], [3], [2]). The simplest terms  $\xi(x, y; *)$  in the first case are

$$ax * y *, \quad xy * a *$$

and in the second one

$$xa * ay **, \quad xx * xy **, \quad xy * yy **.$$

II. The commutative law also satisfies the condition ( $\mathcal{C}$ ). One example of the term  $\xi(x, y; *)$  is:  $xx * xy **$ . This is the well known term which is used to define the ordered pair.<sup>1)</sup>

III. If  $\Sigma(*)$  is a law of the form

$$\Pi x_1 \cdots x_n = \Pi x_{p_1} \cdots x_{p_n}$$

where at both sides of the equality are terms with the same arrangement of operation symbols, and the permutation  $p = \begin{pmatrix} 1 \cdots n \\ p_1 \cdots p_n \end{pmatrix}$  has at least one fixed point, for example  $p(i) = i$ , then one convenient term is

$$\Pi x \cdots xyx \cdots x \quad (y \text{ is at the } i\text{-th place})$$

In the case of laws:

$$xy * z ** = zy * x *, \quad xy * u * vw ** = wy * v * ux **$$

such terms are:  $xy * x *$ ,  $xy * x * xx **$  respectively, for they stay unchanged after applying the corresponding laws [2].

IV. An example of laws of the previous type is

$$\Pi x_1 \cdots x_n = \Pi x_{p_1} \cdots x_{p_n}$$

under the condition that the permutation  $p$  has at least two fixed points, say  $p(i) = i$ ,  $p(j) = j$ . The convenient term  $\xi(x, y; *)$  can be formed as in III, as well as in different way. Namely

$$\Pi a \cdots axa \cdots aya \cdots a \quad (a \text{ is a constant; } x \text{ and } y \text{ are at the } i\text{-th and } j\text{-th place})$$

satisfies also the condition ( $\mathcal{C}$ ).

Remark. If the term  $\xi(x, y; *)$  is such that it depends on the constant  $a$ , i. e. the term of the form  $\xi(x, y, a; *)$ , then  $\text{Def}(\circ)$  becomes:

$$\text{Definition } (\circ_a) \quad xy \circ_a \stackrel{\text{def}}{=} \xi(x, y, a; *)$$

and the operation  $\omega$  can be represented by some term  $t(x_1, \dots, x_n; \circ_\omega)$  i. e. it is not necessary to use a new constant, since  $\circ_\omega$  already depends on  $\omega$ . This situation occurs in the case of associative and entropic laws (if for  $\xi(x, y; *)$  the term  $xa * ay **$  is chosen).

<sup>1)</sup>  $(x, y) \stackrel{\text{def}}{=} \{\{x, x\}, \{x, y\}\} = \{\{x\}, \{x, y\}\}.$

3. The converse of the Theorem 1. is given by the following

**Theorem 2.** *If  $\Sigma(*)$  are algebraic laws such that for each  $\Omega$ -algebra there exists a groupoid  $(G, *)$  satisfying the laws  $\Sigma(*)$ , and  $Q$  is isomorphically embedded in  $(G, *)$ , then there exists a term  $\xi(x, y; *)$  such difficult the operation  $\circ$  defined by  $\text{Def}(\circ)$  does not satisfy any algebraic law.*

**Proof.** If each  $\Omega$ -algebra can be embedded in some groupoid satisfying  $\Sigma(*)$ , then the same holds for  $\circ$ -word algebra  $W$  generated by some set  $X$  and the operation symbol  $\circ$  (of arity two). That is, there exists a groupoid  $(G, *)$  and a term  $\xi(x, y; *)$  such that

$$xy \circ \stackrel{\text{def}}{=} \xi(x, y; *).$$

As  $W$  does not satisfy any algebraic law [1], the term  $\xi(x, y; *)$  is the required term.

4. By analysing the proofs of the previous theorems it is not hard to see that the assumption that  $\Sigma(*)$  are the groupoid laws is not essential. Namely, if  $\Sigma(O)$  are the laws with respect to the operators of some set  $O$ , but such that

*There exists a term  $\xi(x, y; o_1, \dots, o_n)$ ,  $o_i \in O$  such that the operation  $\circ$  (of arity two) defined by*

$$(\mathcal{C}') \quad xy \circ \stackrel{\text{def}}{=} \xi(x, y; o_1, \dots, o_n)$$

*does not satisfy any algebraic law,*

then Theorems 1 and 2 can be extended to the statements about the embedding of each  $\Omega$ -algebra in some  $O$ -algebra satisfying  $\Sigma(O)$ , i.e. belonging to the variety  $V_O(\Sigma)$ . For example, one primitive class  $\Sigma(O)$  satisfying  $(\mathcal{C}')$  is the class  $\Sigma(\{f\})$  with the law:

$$(x_{11}x_{12} \cdots x_{1n}f \cdots x_{n1}x_{n2} \cdots x_{nn}f)f = (x_{11}x_{21} \cdots x_{n1}f \cdots x_{1n}x_{2n} \cdots x_{nn}f)f.$$

The term of  $n$  variables satisfying  $(\mathcal{C}')$  is:

$$(x_{11}a \cdots af \cdots aa \cdots x_{nn}f)f$$

which can easily be reduced to a term of the form  $\xi(x, y, a; f)$ .

The other class  $\Sigma(O)$  is  $\Sigma(\{f, g\})$  ( $f$  of arity one,  $g$  of arity two) with the laws:

$$xxg = xf, \quad xyg = yxg$$

The convenient term  $\xi(x, y; f, g)$  is:  $xfxygg$ . If  $f, g$  are the set operator  $\{\}$  then the term becomes  $\{\{x\}, \{x, y\}\}$ .

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