

SEARCH NUMBER OF THE CARTESIAN PRODUCT OF
GRAPHS

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ABSTRACT

An upper bound for the search number of the Cartesian product $G_1 + G_2$ is determined, where the search numbers of the graphs G_1 and G_2 are $s(G_1)$ and $s(G_2)$, respectively. Using this, some estimates for the search numbers of n -cubes are obtained, for n natural.

Let G be a finite connected graph without loops or multiple edges. We may assume that G is embedded in R^3 so that its vertices v_1, v_2, \dots, v_p are represented by distinct points, and its edges are represented by closed line segments in R^3 which intersect only at the vertices of G . Regarded as a subset of R^3 , G is a topological space with the relative topology. Then G is a compact locally connected metric space in which every connected set is arcwise connected.

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The following definitions are from Parsons ([1], [2]).

Definition 1. A search plan for G is a family $F = \{f_i | 1 \leq i \leq k\}$ of continuous functions $f_i : [0, \infty) \rightarrow G$, such that for every continuous function $g : [0, \infty) \rightarrow G$ there exists $t_0 \in [0, \infty)$ and an $i \in \{1, 2, \dots, k\}$ such that $g(t_0) = f_i(t_0)$.

We think of $g(t)$ and $f_i(t)$ as the positions at time t of a lost man and the i th searcher in the cave represented by graph G , in which the searchers and the lost man move continuously. The searchers must proceed according to a predetermined plan which will capture the lost man even if he were an arbitrarily fast, invisible evader who, clairvoyant, knows the searchers' every move. Then a search plan must catch any possible evader in a finite time.

Definition 2. The search number $s(G)$ of G is the minimum cardinality of all search plans for G .

Our problem is to give an upper bound for the search number of the Cartesian product of given graphs for which the search numbers are known.

Let $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ be two finite connected graphs. We denote by $G_1 + G_2$ the Cartesian product G of the graphs G_1 and G_2 (somewhere called the sum), i.e.

$$G = G_1 + G_2 = (V, E),$$

where

$$V = V_1 \times V_2$$

and

$$E = \{(x_1, y_1)(x_2, y_2) | (x_1 x_2 \in E_1 \text{ and } y_1 = y_2)$$

$$\text{or } (x_1 = x_2 \text{ and } y_1 y_2 \in E_2)\}.$$

Theorem 1. If $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ are two finite connected graphs, then

$$s(G_1 + G_2) \leq \min\{|V_1|s(G_2), |V_2|s(G_1)\} + 1.$$

Proof. Let $V_1 = \{x_1, x_2, \dots, x_p\}$, $V_2 = \{y_1, y_2, \dots, y_k\}$, $s(G_1) = m$, $s(G_2) = n$. Suppose that $pn \leq qm$ and that we have at our disposal $pn + 1$ searchers. Denote these $pn + 1$ searchers by

$$C_0; C_1^1, C_1^2, \dots, C_1^p; C_2^1, C_2^2, \dots, C_2^p; \dots; C_n^1, C_n^2, \dots, C_n^p.$$

Let the vertices $y_{k_1}, y_{k_2}, \dots, y_{k_n}$ (not necessarily different) make a starting position for a search plan in G_2 . (Of course, if $s(G) = n$, for a graph G , then any set of at most n points of G , considered as a subset of R^3 , can serve a starting position for a search plan, but the theorem also holds for digraphs). First, we place all searchers C_j^i ($i \in \{1, 2, \dots, p\}$, $j \in \{1, 2, \dots, n\}$) on the starting vertices in such a way that the searcher C_j^i occupies the vertex (x_i, y_{k_j}) . Now, all the vertices of the copies $G_{y_{k_1}}, G_{y_{k_2}}, \dots, G_{y_{k_n}}$ of the graph G_1 are occupied. (G_{y_k} is the copy of G_1 induced by the vertices $(x_1, y_k), (x_2, y_k), \dots, (x_p, y_k)$ of the graph $G = G_1 + G_2$. Similarly, G_{x_h} is the copy of G_2 induced by the vertices $(x_h, y_1), (x_h, y_2), \dots, (x_h, y_q)$ of the graph $G = G_1 + G_2$). Then the searcher C_0 traverses all the edges of these copies. If the lost man is on some of these edges, he will be found. If not, we begin by simultaneous realizations of p search plans in the copies $G_{x_1}, G_{x_2}, \dots, G_{x_p}$ of G_2 ; the searchers $C_1^1, C_2^1, \dots, C_n^1$ search in the copy G_{x_1} . Each time the searchers C_j^i ($i \in \{1, 2, \dots, p\}$, $j \in \{1, 2, \dots, n\}$) occupy all the vertices of some copies $G_{y_{t_1}}, G_{y_{t_2}}, \dots, G_{y_{t_n}}$ (not necessarily different) of G_1 , they make a pause during which the searcher C_0 traverses all the edges of these copies. The search continues in this way. So, the lost man will be found either in a copy G_{x_i} of G_2 by the searchers $C_1^i, C_2^i, \dots, C_n^i$ or in a copy G_{y_r} of G_2 by the searcher C_0 . \square

Theorem 1 can be generalized in the following way

Theorem 2. Let $G_1 = (V_1, E_1)$, $G_2 = (V_2, E_2), \dots, G_m =$

$= (V_m, E_m)$ be finite connected graphs.

Then

$$s(G_1 + G_2 + \dots + G_m) \leq \min_{1 \leq i \leq m} \{s(G_i) \prod_{\substack{1 \leq j \leq m \\ j \neq i}} |V_j|\} + 1.$$

Proof. Similar as for Theorem 1. \square

From Theorem 2, we obtain an upper bound for the search number of the n -dimensional cube Q_n .

Corollary 1.

$$s(Q_n) \leq 2^{n-1} + 1.$$

For $n = 1$ and $n = 2$, the strict inequality holds. Namely, $s(Q_1) = 1$ and $s(Q_2) = 2$.

For $n = 3$, the equality holds, i.e. $s(Q_3) = 5$, but a rigorous proof requires some care.

Similarly, for the search number of the Cartesian product of paths we obtain:

Corollary 2.

$$s(P_{n_1} + P_{n_2} + \dots + P_{n_m}) \leq \frac{n_1 n_2 \dots n_m}{\max\{n_1, n_2, \dots, n_m\}} + 1.$$

REFERENCES

- [1] T.D. Parsons: Pursuit-evasion in a graph, Lecture Notes in Math., 642, Springer, Berlin, (1978) 426 - 441.
- [2] T.D. Parsons: The search number of a connected graph Proc. 9th South-Eastern Conf. on Combinatorics, Graph Theory and Computing, (1978), 549 - 554.

REZIME

ISTRAŽNI BROJ KARTEZIJEVOG PROIZVODA GRAFOVA

U radu je određena jedna gornja granica za s-broj Dekartovog proizvoda $G_1 + G_2$, za poznate s-brojeve $s(G_1)$ i $s(G_2)$. Kao posledica dobijena je procena za broj $s(Q_n)$, n je prirodan broj, gde je Q_n n -dimenzionalna kocka.

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