

NOTE ON THE STANDARD MATRIX REPRESENTATION
OF A MATROID

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ABSTRACT

The main result of this paper is: If $[1_k, A]$ is the SMR of a matroid with respect to a base B , then $\text{rank}(A) \geq C(B)$, where $C(B)$ denotes the maximal number of pairwise disjoint fundamental circuits with respect to B .

INTRODUCTION.

The matroid theory terminology and results used in this paper conform to standard literature (see for example [3, 4]).

Let E be a finite set, and $M := M(E, r)$ a matroid on E with r as the rank function ($r: 2^E \rightarrow \mathbb{N}$, where 2^E is the power set of E , and \mathbb{N} the set of non-negative integers). A subset $S \subseteq E$ is called independent if $r(S) = |S|$, where $|S|$ denotes the cardinality of S . Let $F(M)$ be the family of independent sets of M . A basis of M is a maximal independent subset of E . A subset of E which is not independent is called dependent, and a minimal dependent subset of E is a

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circuit. We shall denote by $C(M)$ the family of circuits of M . A circuit of cardinality 1 is a loop of M . For every $S \subseteq E$, \bar{S} denotes the span of S , i.e. $\bar{S} = \{e \in E : r(S \cup \{e\}) = r(S)\}$.

If M is a matroid with $k = r(E)$, let $B = \{e_1, e_2, \dots, e_k\}$ be a basis of M and $E - B = \{f_1, f_2, \dots, f_m\}$. If M is representable over a field \mathbb{F} (see [1]), it will have a standard matrix representation (see also [2]), with respect to the basis B , of the form $R = [I_k, A]$, where I_k is the identity matrix of order k and A is a $k \times m$ matrix with entries belonging to \mathbb{F} . If $E_1, E_2, \dots, E_k, F_1, F_2, \dots, F_m$ denote the column vectors of R , and the map σ , defined by $\sigma(e_i) = E_i, i=1, 2, \dots, k$ and $\sigma(f_i) = F_i, i=1, 2, \dots, m$, is a representation of M over \mathbb{F} , in the sense that a subset S of E belongs to $F(M)$ if and only if the corresponding vectors of $\sigma(S)$ are linearly independent over \mathbb{F} (see [3]).

If B is any basis of M and $\{f_1, f_2, \dots, f_m\} = E - B$, then there exists a unique $C_i \in C(M)$, containing f_i and, otherwise, elements of B only. This circuit C_i is called fundamental with respect to f_i and B and will be denoted by $C(f_i, B)$. For a fixed basis B of a representable matroid M , we denote by $c(B)$ the maximal number of pairwise disjoint fundamental circuits with respect to $E - B$ and B . The main result of this paper is the following:

THEOREM. *If $R = [I_k, A]$ is the standard matrix representation with respect to B , then $\text{rank}(A) \geq c(B)$.*

The key lemma. Let B be a fixed basis of M and $e \in E$ arbitrarily chosen. Obviously, $e \in \bar{B}$. Then we can consider the family of sets $D(e, B) = \{S \subseteq B : e \in \bar{S}\}$.

We shall denote by $B(e)$ a minimal (with respect to the inclusion of sets) element of $D(e, B)$. It can be checked that

$$\begin{aligned} \text{if } r(\{e\}) = 0, & \text{ then } B(e) = \emptyset, \\ \text{if } e \in B, & \text{ then } B(e) = \{e\}. \end{aligned}$$

Suppose that $e \notin B$, and let $C(e, B)$ be the fundamental circuit with respect to e and B . Obviously, by definition, $B(e) \cup \{e\}$ is a circuit of M contained in $B \cup \{e\}$, and therefore $B(e) \cup \{e\} = C(e, B)$. Thus, $B(e)$ is uniquely defined. Throughout, we shall make use of the following fundamental properties of the matroids:

- (P1) If $C_1, C_2 \in C(M)$ such that $C_1 \subseteq C_2$, then $C_1 = C_2$.
- (P2) If C_1, C_2 are two distinct circuits of M such that $e \in C_1 \cap C_2$, then there exists $C_{12} \in C(M)$ with $C_{12} \subseteq (C_1 \cup C_2) - \{e\}$.
- (P3) A subset of E is independent if and only if it does not contain a circuit.
- (P4) A subset of E is independent if and only if it is contained in a basis.

LEMMA. Let B be a fixed basis of M and $F \subseteq E$. If F does not contain any loop and $B(e) \cap B(f) = \emptyset$ for every distinct $e, f \in F$, then $F \in \mathcal{F}(M)$.

P r o o f. Let $F = \{e_1, e_2, \dots, e_t\}$. If $F \subseteq B$, then the theorem is trivial, by (P4). Thus, two cases must be considered:

- (a) $F \cap B = \emptyset$,
- (b) $F = H \cup G$, $H \cap B = \emptyset$, $G \subseteq B$ and $G \neq \emptyset$.

We prove now the lemma in case (a). Let us denote $B(e_1) = B_1$. We have seen above that $B_1 \cup \{e_1\} = C(e_1, B)$.

Suppose that $F \notin \mathcal{F}(M)$. By (P3), there exists $C \in C(M)$ such that $C \subseteq F$. Without loss of generality, we can consider C of the form $C = \{e_1, e_2, \dots, e_s\}$. By (P2), there exists a circuit $C_1 \subseteq (C \cup C(e_1, B)) - \{e_1\}$. If $C_1 \subseteq C(e_1, B)$, then $C_1 = C(e_1, B)$ by (P1), and this is impossible. Thus, $C_1 \subseteq (B_1 \cup C) - \{e_1\}$ and $C_1 \neq C$. Similarly, there exists $C_2 \in C(M)$ with $C_2 \subseteq (B_2 \cup C) - \{e_2\}$ and $C_2 \neq C$. Suppose that $e_1 \notin C_2$. By (P2), there exists $\tilde{C} \in C(M)$

such that $\tilde{C} \subseteq (C \cup C_2) - \{e_1\}$, $\tilde{C} \neq C_2$. Since $B_1 \cap B_2 = \emptyset$ (by hypothesis), we then have $\tilde{C} \neq C_1$. If $e_2 \in \tilde{C} \cap C_1$, then, by (P2), there exists $\tilde{\tilde{C}} \in C(M)$ with $\tilde{\tilde{C}} \subseteq (\tilde{C} \cup C_1) - \{e_1, e_2\}$. Therefore, there exists $C_0 \in C(M)$ such that $C_0 \subseteq (B_1 \cup B_2 \cup C) - \{e_1, e_2\}$.

On the other hand, $C_3 = B_3 \cup \{e_3\}$ is a circuit of M and $C_3 \neq C_0$. If $e_3 \in C_0$, then, by (P2), there exists $C'_0 \in C(M)$ with $C'_0 \subseteq (C_0 \cup C_3) - \{e_1, e_2, e_3\}$. Thus, there exists C_0 (if $e_3 \in C_0$, we take $C_0 := C'_0$) in $C(M)$ such that $C_0 \subseteq (B_1 \cup B_2 \cup B_3 \cup C) - \{e_1, e_2, e_3\}$. Repeating the above a finite number of times ($s-3$ times), thus, gives rise to a circuit C_0 with $C_0 \subseteq B_1 \cup B_2 \cup \dots \cup B_s$, i.e. a subset of B contains a circuit, contradicting (P3) and (P4). Hence, $F \in F(M)$.

Now, we shall prove the lemma in case (b). Suppose that $F \notin F(M)$, i.e. by (P3), there exists $C \in C(M)$ such that $C \subseteq F$. Obviously, C must contain at least an element of H as otherwise $G \notin F(M)$, which is in contradiction with (P4). Considering C of the form $C = \{e_1, e_2, \dots, e_s\} \cup T$ with $\{e_1, e_2, \dots, e_s\} \subseteq H$ and $T \subseteq G$, and repeating the above judgement (as in case (a)) for the set $\{e_1, e_2, \dots, e_s\}$, we obtain a similar contradiction: a subset of B contains a circuit. Therefore, the lemma is entirely proved. (QED).

COROLLARY. *Let B be a fixed basis of M , and $\{f_{i_1}, f_{i_2}, \dots, f_{i_t}\} \subseteq E - B$ such that $C(f_{i_j}, B)$ are pairwise disjoint. Then $\{f_{i_1}, f_{i_2}, \dots, f_{i_t}\} \in F(M)$.*

P r o o f. It readily follows from the lemma (QED).

Proof of the theorem. Let M be a representable matroid over the field F , B a fixed basis of M and $R = [I_k, A]$ the standard matrix representation with respect to B . By the above corollary, it follows that the column vectors

$F_{i_1}, F_{i_2}, \dots, F_{i_t}$, are linearly independent over F , i.e.

$\text{rank}(A) \geq t$ (QED).

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REZIME

NOTA O STANDARDNOJ MATRIČNOJ REPREZENTACIJI MATROIDA

Osnovni rezultat ovog rada je: Ako je $[I_k, A]$ SMR matroida u odnosu na bazu B , tada je $\text{rank}(A) \geq C(B)$, gde je $C(B)$ maksimalan broj.