

AN APPLICATION OF PSEUDO-BOOLEAN FUNCTIONS TO THE TREE THEORY

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Let $L = \{a_0, a_1, \dots, a_p\}$ with $p \geq 1$, $p \in N$, L^n be the Cartesian product of L ; let $(P, +, \cdot)$ be a commutative ring with a unit element 1; with no zero divisors; and let the order on P be defined by a binary relation \geqslant .

A function $f: L^n \rightarrow P$ is called a pseudo-Boolean function.

Let us define a function from L^2 to P , x_i^b :

$$(1) \quad \begin{aligned} x_i^b &= 1 \text{ iff } (x_i, b) = (x, x) \\ x_i^b &= 0 \text{ in the other case } (i=1, \dots, n), \end{aligned}$$

where 0 is the zero, 1 the unit element of P . Directly from (1) follows:

$$(2) \quad \sum_{b \in L} x_i^b = 1, \quad i=1, \dots, n.$$

Each generalised pseudo-Boolean function f can be written

$$(3) \quad \begin{aligned} f(x_1, \dots, x_n) &= \sum f(b_1, \dots, b_n) \cdot x_1^{b_1} \cdots x_n^{b_n}, \\ (b_1, \dots, b_n) &\in L^n \end{aligned}$$

it is evident in [2].

If we take in (2) a fixed point $C = (c_1, c_2, \dots, c_n)$ from L^n , we get

$$(4) \quad x_i^{c_i} = 1 - \sum_{b_i \in L \setminus \{c_i\}} x_i^{b_i}, \quad i=1, \dots, n,$$

where $a + (-a) = 0$, $a \in P$.

If we put (4) and (2) in (3) we can transform the generalised pseudo-Boolean function f in the form

$$(5) \quad f(x_1, \dots, x_n) = d + \sum_{j=t_1}^{t_n} \sum_{b \in L \setminus \{c_{i_j}\}} d_{jb} x_j^b, \quad d = f(c)$$

where i_1, i_2, \dots, i_n is the permutation of the set $\{1, 2, \dots, n\}$, $|j|$ is the index of j .

The generalised pseudo-Boolean functions

$$(6) \quad d_{jb} = d_{jb}(x_{|j|+1}, \dots, x_n), \quad j=i_1, \dots, i_n \quad b \in L \setminus \{c_{|j|}\}$$

$$(6') \quad d_{i_1 c_1}(x) = \dots = d_{i_n c_n}(x) = 0, \quad x = (x_1, \dots, x_n),$$

are called the derivatives of f .

According to (5) for $L = \{0, 1, \dots, p\}$ and $C = \{p, \dots, p\}$ the generalised pseudo-Boolean function f can be transformed in

$$(7) \quad f(x) = d + \sum_{i=1}^n \sum_{\alpha \in L'} d_{i\alpha}(x_{i+1}, \dots, x_n) x_i^\alpha + \sum_{\alpha \in L'} d_{n\alpha} x_n^\alpha,$$

$$(L' = L \setminus \{p\}),$$

where $d = f(p, \dots, p)$; the derivatives d_{ij} , $i=1, \dots, n-1$; $j \in L'$ are generalised pseudo-Boolean functions

$$(8) \quad d_{ij}: L^{n-i} \rightarrow P, \quad i=1, \dots, n-1, \quad j \in \{0, 1, \dots, p-1\},$$

while the derivatives

$$(8') \quad d_{n0}, d_{n1}, \dots, d_{np-1}$$

are constants in P .

THEOREM 1. One and only one oriented tree (q_n, X_n) with the root d , with n levels and $(p+1)^{n-1} - 1$ poles, where $X_n = \{f(x_1, \dots, x_n) \mid (x_1, \dots, x_n) \in L^n\}$,

$$f(p, \dots, p, (\alpha)_i, p, \dots, p) = f(\underbrace{p, \dots, p}_{i-1\text{-times}}, \alpha, p, \dots, p),$$

$$(i) \quad q_n(f(p, \dots, p)) = f(p, \dots, p, (\alpha_1)_{i_1}, p, \dots, p), \\ i_1 = 1, 2, \dots, n; \quad \alpha_1 \in L',$$

$$(ii) \quad q_n(f(p, \dots, p, (\alpha_1)_{i_1}, p, \dots, p)) = f(p, \dots, p, (\alpha_2)_{i_2}, p, \dots, p, (\alpha_1)_{i_1}, p, \dots, p), \\ i_1 = 2, 3, \dots, n; \quad i_2 = 1, 2, \dots, n-1; \\ i_2 < i_1; \quad \alpha_1, \alpha_2 \in L',$$

$$(iii) \quad q_n(f(p, \dots, p, (\alpha_2)_{i_2}, p, \dots, p, (\alpha_1)_{i_1}, p, \dots, p)) = \\ = f(p, \dots, p, (\alpha_3)_{i_3}, p, \dots, p, (\alpha_2)_{i_2}, p, \dots, p, (\alpha_1)_{i_1}, p, \dots, p) \\ i_1 = 3, 4, \dots, n; \quad i_2 = 2, 3, \dots, n-1; \quad i_3 = 1, 2, \dots, n-1; \\ i_3 < i_2 < i_1; \quad \alpha_1, \alpha_2, \alpha_3 \in L',$$

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$$(n=1 \text{ i}) \quad q_n(f(p, \dots, p, (\alpha_{n-1})_{i_{n-1}}, p, \dots, p, (\alpha_1)_{i_1}, p)) = f(\alpha_n, \alpha_{n-1}, \dots, \alpha_1) \\ i_1 = n-1, n; \quad i_2 = n-2, n-1; \dots; i_{n-1} = 1, 2; \quad i_{n-1} < i_{n-2} < \dots < i_2 < i_1; \\ \alpha_1, \alpha_2, \dots, \alpha_n \in L',$$

can be associated to each injective generalised pseudo-Boolean function f .

Proof. If $n=1$ the generalised pseudo-Boolean function f has the following form

$$f(x_1) = d + \sum_{\alpha \in L'} f d_{1\alpha} x_1^\alpha, \quad x_1 \in L,$$

where $f(p)=d$, $f(\alpha)=d+d_{1\alpha}$, $\alpha \in L'$.

The oriented tree (q_1, X_1) , associated to the function $f(x_1)$ has a root d and one level with p poles (fig. 1.),

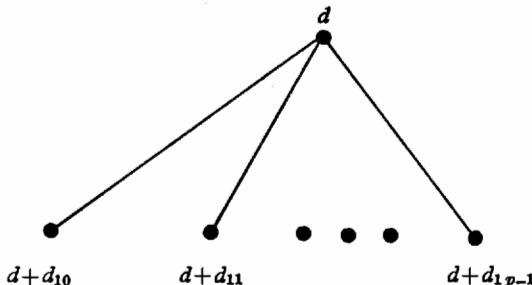


Fig. 1

where

$$q_1(d) = d + d_{1\alpha}, \quad \alpha \in L'$$

$$X_1 = \{f(x_1) \mid x_1 \in L\}.$$

If $n=2$, the generalised pseudo-Boolean function f has the following form

$$f(x_1, x_2) = d + \sum_{\alpha \in L'} d_{1\alpha}(x_2) x_1^\alpha + \sum_{\alpha \in L'} d_{2\alpha} x_2^\alpha, \quad (x_1, x_2) \in L^2.$$

The oriented tree (q_2, x_2) associated to the function $f(x_1, x_2)$ (Fig. 2) has:

a) A root $f(p, p)=d$.

b) Poles of the first level

$$f(\alpha, p) = d + d_{1\alpha}(p), \quad \alpha \in L'$$

$$f(p, \alpha) = d + d_{2\alpha}, \quad \alpha \in L'.$$

c) Poles of the second level

$$f(\alpha, \beta) = d + d_{2\beta} + d_{1\alpha}(\beta), \quad \alpha, \beta \in L'.$$

According to a), b) and c) the oriented tree (q_2, X_2) associated to the function $f(x_1, x_2)$ has a root d , two levels with $(p+1)^2 - 1$ poles (fig. 2.), where

$$b') q_2(p) = d + d_{1\alpha}(p), \quad \alpha \in L'$$

$$q_2(p) = d + d_{2\alpha}, \quad \alpha \in L'.$$

$$c') q_2(d + d_{2\alpha}) = d + d_{2\beta} + d_{1\alpha}(\beta), \quad \alpha, \beta \in L'.$$

We conclude by the induction that one and only one oriented tree (q_n, x_n) with a root d , with n levels and with $(p+1)^n - 1$ poles can be associated to the injective generalised pseudo-Boolean function of the form (7).

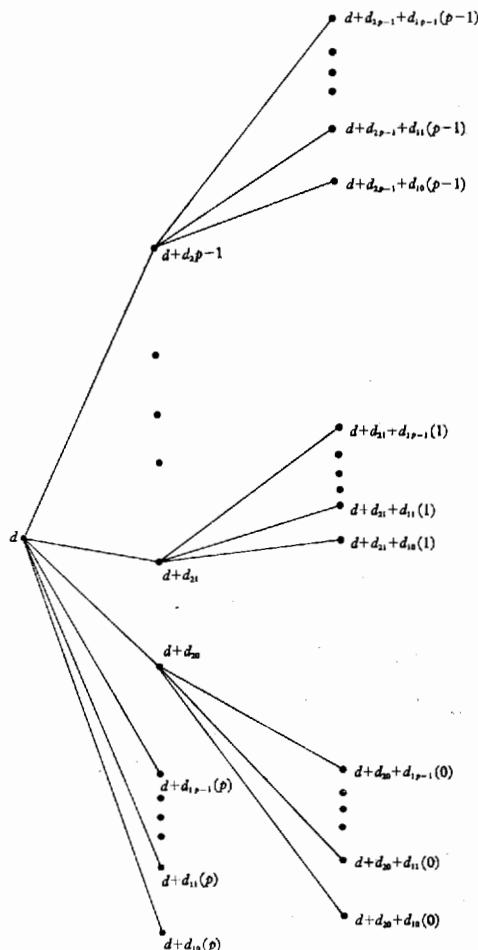


Fig. 2

THEOREM 2. If the derivatives (8) and (8') are $d_{ij} > 0$ ($d_{ij} < 0$) $i = 1, \dots, n$, $j \in \{0, 1, \dots, p-1\}$ then the generalised pseudo-Boolean function (7) has a unique minimum at (p, p, \dots, p) the minimum is $f_{\min}(p, \dots, p) = d$ 1 (a unique maximum $f_{\max}(p, \dots, p) = d$).

Proof. Let the derivatives (8) and (8') be positive (negative); $d_{ij} > 0$ ($d_{ij} < 0$) $i = 1, \dots, n$, $j \in \{0, 1, \dots, p-1\}$. For every vector $(\alpha_1, \dots, \alpha_n)$ from L^n function (7) has the form

$$(9) \quad f(\alpha_1, \dots, \alpha_n) - f(p, \dots, p) = d_{1\alpha_1}(\alpha_2, \dots, \alpha_n) + d_{2\alpha_2}(\alpha_3, \dots, \alpha_n) + \dots + d_{n-1\alpha_{n-1}}(\alpha_n) + d_{n\alpha_n}.$$

As we supposed that the derivatives $d_{ij} > 0$ ($d_{ij} < 0$), from (9) follows that for every $(x_1, \dots, x_n) \in L^n$

$$f(x_1, \dots, x_n) - f(p, \dots, p) > 0 \quad (f(x_1, \dots, x_n) - f(p, \dots, p) < 0).$$

Let us prove that the function $f(7)$ has a unique minimum. Let us suppose that the function $f(7)$ has two points of minimum (maximum).

$$(10) \quad (c_1, \dots, c_n) \neq (p, \dots, p) \text{ and } f_{\min}(c_1, \dots, c_n) = f_{\min}(p, \dots, p) \\ (f_{\max}(c_1, \dots, c_n) = f_{\max}(p, \dots, p)).$$

If we put (c_1, \dots, c_n) in (7) we get (11). We supposed that $d_i c_i > 0$, $i=1, \dots, n-1$, $d_n > 0$ ($d_i c_i < 0$, $d_n < 0$, $i=1, \dots, n-1$).

This implies

$$f_{\min}(c_1, \dots, c_n) \approx f_{\min}(p, \dots, p) = d.$$

$$(11) \quad f(c_1, \dots, c_n) = f(p, \dots, p) + d_1 c_1 (c_2, \dots, c_n) + d_2 c_2 (c_3, \dots, c_n) + \\ + \dots + d_{n-1} c_{n-1} (c_n) + d_{c_n}.$$

It is in contradiction with (10).

Theorem 2 is proved.

COROLLARY. Every injective generalised pseudo-Boolean function $f: L^n \rightarrow P$ is associated with $n!$ oriented trees with a root d , n levels and $(p+1)^n - 1$ poles, the poles are toe values of the function f .

Proof. According to (5), a generalised pseudo-Boolean function can be transformed in $n!$ forms and according to Theorem 1 we have $n!$ oriented trees.

EXAMPLE. Let the function $f: \{0, 1, 2\}^2 \rightarrow R$ (R is a set of real numbers) be written in the following way

$$f(x, y) = x^0 d_{10}(y) + d_{11}(y)x^1 + d_{20}y^0 + d_{21}y^1 + f(2, 2)$$

The corresponding oriented tree is in Fig. 3.

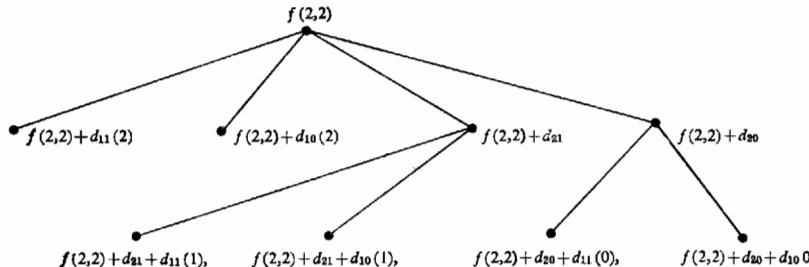


Fig. 3

The other problem is to associate the given tree (q, x) with a generalised pseudo-Boolean function. Let an oriented tree be given in Fig. 4.

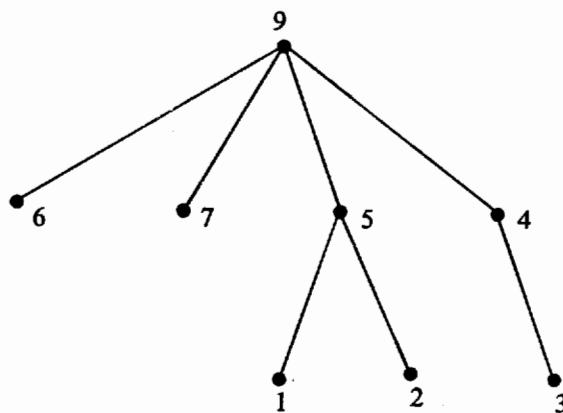


Fig. 4

Let us compare the trees from Fig. 3 and 4. Another pole can be added to the tree in Fig. 5.

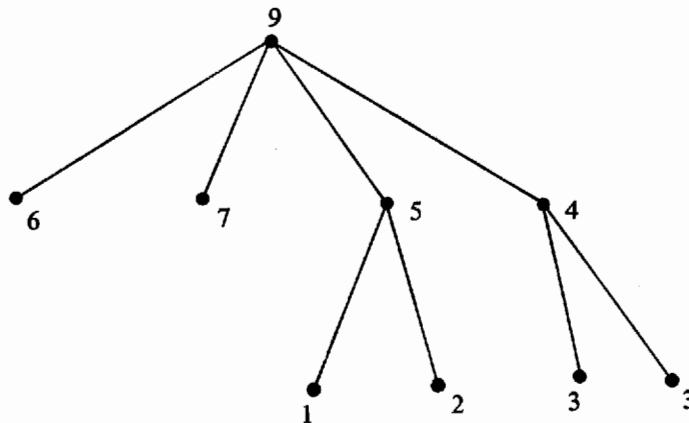


Fig. 5

Let us compare the trees from Fig. 5 and Fig. 3 and let us find the derivatives d_{21} , d_{20} , $d_{10}(y)$, $d_{11}(y)$ when $d_{20} = -5$, $d_{21} = -4$, and $d_{10}(y)$, $d_{11}(y)$ are given

y	0	1	2
$d_{10}(y)$	-1	-3	-2

y	0	1	2
$d_{11}(y)$	-1	-4	-3

where

$$d_{10}(y) = -y^{\circ} - 3y^1 - 2y^2$$

$$d_{11}(y) = -y^{\circ} - 4y^1 - 3y^2.$$

Therefore the generalised function

$$f(x, y) = x^{\circ}(-y^{\circ} - 3y^1 - 2y^2) + x^1 \cdot (-y^{\circ} - 4y^1 - 3y^2) - 5y^{\circ} - 4y^1 + 9$$

is associated to the tree given by Fig. 5.

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PRIMENA GENERALISANIH PSEUDO-BULOVIH FUNKCIJA U TEORIJI STABALA

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REZIME

Dokazano je da se svakoj injektivnoj generalisanoj pseudo-Bulovoj funkciji može pridružiti jedno jedino orijentisano stablo. Takođe je dokazano da generalisana pseudo-Bulova funkcija sa pozitivnim (negativnim) derivatima ima jedinstvenim minimum (maksimum).