

## HOSOYA POLYNOMIAL OF HANOI GRAPHS

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ABSTRACT. A recursive method for the calculation of the Hosoya polynomial of Hanoi graph is designed. This make it possible to compute various distance-based invariants of Hanoi graphs.

### 1. INTRODUCTION

Let  $G$  be a connected graph and  $V(G) = \{v_1, v_2, \dots, v_p\}$  its vertex set. The distance of the vertices  $v_i$  and  $v_j$ , denoted by  $d(v_i, v_j|G)$ , is the length of (= number of edges in) a shortest path that connects  $v_i$  and  $v_j$  [2]. If  $i = j$ , then  $d(v_i, v_j|G) = 0$ .

A large number of distance-based graph invariants has been studied both in the last 50–60 years and in the recent past (see [4, 5, 3, 1] and the references cited therein). Of these we mention here the Wiener index ( $W$ ), the hyper-Wiener index ( $WW$ ), the Harary index ( $Ha$ ), the reciprocal Wiener index ( $RW$ ), and the distance moments ( $W_\rho$ ). These are defined as

$$\begin{aligned}W &= W(G) &= \sum_{i < j} d(v_i, v_j|G), \\WW &= WW(G) &= \frac{1}{2} \sum_{i < j} \left[ d(v_i, v_j|G)^2 + d(v_i, v_j) \right], \\Ha &= Ha(G) &= \sum_{i < j} \frac{1}{d(v_i, v_j|G)^2}, \\RW &= RW(G) &= \sum_{i < j} \frac{1}{d(v_i, v_j|G)}, \\W_\rho &= W_\rho(G) &= \sum_{i < j} d(v_i, v_j|G)^\rho.\end{aligned}$$

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In the study of these and other distance-based graph invariants it is often convenient to work with the numbers  $d(G, k)$  which are equal to the number of pairs of vertices of the graph  $G$  that are at distance  $k$ . With this notation, we have

$$(1.1) \quad W = W(G) = \sum_{k \geq 1} k d(G, k),$$

$$(1.2) \quad WW = WW(G) = \frac{1}{2} \sum_{k \geq 1} (k^2 + k) d(G, k),$$

$$(1.3) \quad Ha = Ha(G) = \sum_{k \geq 1} \frac{1}{k^2} d(G, k),$$

$$(1.4) \quad RW = RW(G) = \sum_{k \geq 1} \frac{1}{k} d(G, k),$$

$$(1.5) \quad W_\rho = W_\rho(G) = \sum_{k \geq 1} d(G, k) k^\rho.$$

One of the advantages of the latter formulas is that these can be applied also in the case of graphs that are not connected.

In the 1980s Hosoya [8] conceived a graph polynomial which he named "Wiener polynomial", but which most contemporary authors call the "Hosoya polynomial". It is defined as

$$(1.6) \quad H(G; \lambda) = \sum_{k \geq 1} d(G, k) \lambda^k.$$

The maximal distance between two vertices of the graph  $G$  is referred to as the *diameter* of  $G$ . It will be denoted by  $d(G)$ . Then  $d(G, k) > 0$  holds for  $1 \leq k \leq d(G)$ , whereas  $d(G, k) = 0$  holds for  $k > d(G)$ . In view of this, we see that  $H(G; \lambda)$  is a polynomial of degree  $d(G)$ .

The theory of the Hosoya polynomial is nowadays well elaborated; its details and additional references can be found in the recent survey [6]. By comparing Eq. (1.6) with Eqs. (1.1)–(1.5), it should be evident that if one knows the Hosoya polynomial of a graph, then it is elementary to compute any of its distance-based invariants. The respective formulas can be found elsewhere [6]. The simplest such formula is

$$W(G) = H'(G; 1)$$

where  $H'(G; \lambda)$  denotes the first derivative of  $H(G; \lambda)$ . This result motivated Hosoya to call  $H(G; \lambda)$  the Wiener polynomial [8].

## 2. THE HANOI GRAPH

The tower of Hanoi puzzle, invented in 1883 by the French mathematician, Edouard Lucas, has become a classic example in the analysis of algorithms and discrete mathematics [9, 7].

The puzzle consists of  $n$  discs, no two of the same size, stacked on three vertical pegs, in such a way that no disc lies on top of a smaller disc. A permissible move

is to take a top disc from one of the pegs and move it to one of the other pegs as long as it is not placed on top of a smaller disc. The set of configurations of the puzzle, together with the permissible moves, thus forms a graph in a natural way. The number of vertices in the  $n$ -disc Hanoi graph  $H_n$  is  $3^n$  and the number of edges is  $3(3^n - 1)/2$ . One can verify that each Hanoi graph has a unique Hamiltonian cycle and  $3^{n-1}$  small triangles.

The following figure shows the Hanoi graphs for  $n = 1, 2, 3$ .

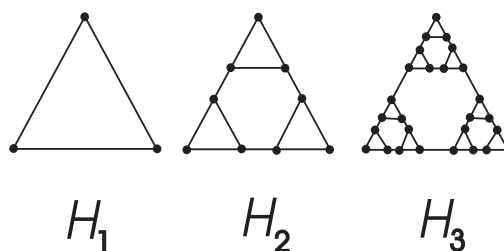


FIGURE 1. The first three Hanoi graphs.

The Hanoi graph  $H_n$  can be constructed by taking the vertices to be the odd binomial coefficients of the Pascal triangle, computed on the integers from 0 to  $2^n - 1$  and drawing an edge whenever the coefficients are adjacent diagonally or horizontally [9]. More information on shortest paths in the tower of Hanoi graph and finite automata can be found in [10].

The main properties of Hanoi graphs are summarized in the following proposition:

**Proposition 2.1.** [9] *For all positive integers  $n$ ,  $H_n$  is a planar 2-connected Hamiltonian graph of order  $3^n$  and diameter  $d = 2^n - 1$ , with exactly three vertices of degree 2 and all other vertices of degree 3.*

In the subsequent section we show how the Hosoya polynomial of the Hanoi graphs can be recursively evaluated.

### 3. HOSOYA POLYNOMIAL OF HANOI GRAPHS

The Hosoya polynomial of the Hanoi graph  $H_n$  is of the form

$$H(H_n; \lambda) = d(H_n, 1) \lambda + d(H_n, 2) \lambda^2 + \cdots + d(H_n, d) \lambda^d$$

where  $d = 2^n - 1$  is the diameter of  $H_n$ .

**Theorem 3.1.** *Let  $n \geq 1$  and let the ordered  $(2^{n-1})$ -tuple of integers  $S_n$  be recursively defined as:*

$$\begin{aligned} (3.1) \quad S_1 &= \{1\}, \\ S_2 &= \{1, 2\} = \{S_1, 2S_1\}, \\ (3.2) \quad S_3 &= \{1, 2, 2, 4\} = \{S_2, 2S_2\}, \end{aligned}$$

$$(3.3) \quad S_4 = \{1, 2, 2, 4, 2, 4, 4, 8\} = \{S_3, 2S_3\},$$

...

$$(3.4) \quad S_n = \{S_{n-1}, 2S_{n-1}\}.$$

Denote the elements of  $S_n$  by  $\ell_i$ ,  $i = 0, 1, 2, \dots, 2^{n-1} - 1$  so that

$$S_n = \{\ell_0, \ell_1, \ell_2, \dots, \ell_{2^{n-1}-1}\}$$

and assume that  $\ell_t = 0$  for  $t < 0$  and  $t > 2^{n-1} - 1$ . For  $k = 1, 2, \dots, 2^n - 1$  define,

$$(3.5) \quad b(n, k) = \sum_{i=0}^k \ell_i \ell_{k-i-1}.$$

Then the coefficients of the Hosoya polynomial of the Hanoi graph  $H_n$  satisfy the recursion relation:

$$(3.6) \quad d(H_n, k) = 3[d(H_{n-1}, k) + b(n, k)]$$

for  $k = 1, 2, \dots, 2^n - 1$ .

*Proof.* Let  $H_n$  be the Hanoi graph with the vertex set  $V(H_n)$  partitioned into disjoint subsets  $A$ ,  $B$ , and  $C$  as shown in Figure 2

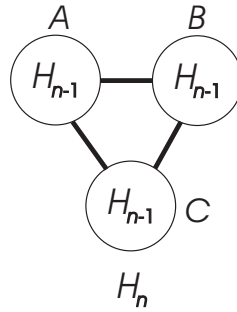


FIGURE 2. Partitioning of the vertices of the Hanoi graph  $H_n$ .

For  $X, Y = A, B, C$ , denote by  $d_{XY}(H_n, k)$  the number of vertex pairs  $(v_i, v_j)$  of the Hanoi graph  $H_n$  that are at distance  $k$ , so that  $v_i \in X$  and  $v_j \in Y$ . Then,

$$(3.7) \quad \begin{aligned} d(H_n, k) &= d_{AA}(H_n, k) + d_{BB}(H_n, k) + d_{CC}(H_n, k) \\ &+ d_{AB}(H_n, k) + d_{AC}(H_n, k) + d_{BC}(H_n, k). \end{aligned}$$

In view of the notation described in Figure 2,

$$(3.8) \quad d_{AA}(H_n, k) = d_{BB}(H_n, k) = d_{CC}(H_n, k) = d(H_{n-1}, k).$$

In view of the symmetry of the Hanoi graphs (cf. Figures 1 and 2),

$$d_{AB}(H_n, k) = d_{AC}(H_n, k) = d_{BC}(H_n, k).$$

Therefore, we only need to compute  $d_{AB}(H_n, k)$ .

Let  $G_{AB}$  be the subgraph induced by the sets  $A$  and  $B$  as shown in Figure 3.

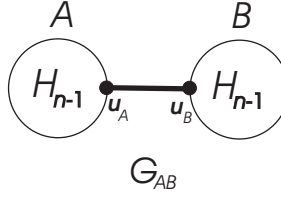


FIGURE 3. The subgraph induced by the vertex sets  $A$  and  $B$  of the Hanoi graph  $H_n$ .

Evidently,  $d_{AB}(H_n, k)$  is equal to the number of vertex pairs of  $G_{AB}$  whose distance is  $k$ , such that one vertex belongs to  $A$  and the other to  $B$ . From Fig. 3 we see that the number of such vertex pairs is equal to the sum over  $i = 0, 1, \dots, k$ , of the product of the number of vertices in  $A$  whose distance to vertex  $u_A$  is  $i$  and the number of vertices in  $B$  whose distance to vertex  $u_B$  is  $k - i - 1$ . If these numbers are denoted by  $\ell_i$  and  $\ell_{k-i-1}$ , then we have

$$(3.9) \quad d_{AB}(H_n, k) = \sum_{i=0}^k \ell_i \ell_{k-i-1} \equiv b(n, k)$$

as well as

$$(3.10) \quad d_{AC}(H_n, k) = d_{BC}(H_n, k) = b(n, k).$$

In order to complete the proof of Theorem 3.1, it only remains to verify that  $\ell_i$  counts the number of vertices in  $H_{n-1}$  whose distance to vertex  $u_A$  (or more generally, to a vertex of degree 2) is equal to  $i$ .

That this indeed is the case can be directly checked for  $H_1$ ,  $H_2$ , and  $H_3$ , resulting in the sequences (3.1), (3.2), and (3.3), respectively. That the general form of such a sequence is (3.4) is now easily proved by induction on  $n$ .

Thus Eq. (3.9) is valid, and therefore also Eq. (3.10) holds. Substituting Eqs. (3.9), (3.10), and (3.8) back into Eq. (3.7), we arrive at the recursive formula (3.6), by which the proof of Theorem 3.1 is completed.  $\square$

#### 4. EXAMPLES

As illustrative examples of the above result, we compute  $H(H_2; \lambda)$  and  $H(H_3; \lambda)$ , which are polynomials in  $\lambda$  of degree  $2^2 - 1 = 3$  and  $2^3 - 1 = 7$ , respectively.

As easily seen,  $H(H_1; \lambda) = 3\lambda$ , i. e.,  $d(H_1, 1) = 3$  and  $d(H_1, k) = 0$  for  $k \geq 2$ .

We first need  $S_2 = \{1, 2\}$ , which means that

$$\begin{aligned} \ell_0 &= 1 \\ \ell_1 &= 2 \end{aligned}$$

Now, by using formula (3.5),

$$\begin{aligned} b(2, 1) &= \ell_0 \ell_0 = 1 \cdot 1 = 1, \\ b(2, 2) &= \ell_0 \ell_1 + \ell_1 \ell_0 = 1 \cdot 2 + 2 \cdot 1 = 4, \\ b(2, 3) &= \ell_0 \ell_2 + \ell_1 \ell_1 + \ell_2 \ell_0 = 1 \cdot 0 + 2 \cdot 2 + 0 \cdot 1 = 4. \end{aligned}$$

From Eq. (3.6) we now get

$$\begin{aligned} d(H_2, 1) &= 3[d(H_1, 1) + b(2, 1)] = 3[3 + 1] = 12, \\ d(H_2, 2) &= 3[d(H_1, 2) + b(2, 2)] = 3[0 + 4] = 12, \\ d(H_2, 3) &= 3[d(H_1, 3) + b(2, 3)] = 3[0 + 4] = 12 \end{aligned}$$

which finally gives

$$H(H_2; \lambda) = 12\lambda + 12\lambda^2 + 12\lambda^3.$$

Repeating the calculation for  $n = 3$ , we start with  $S_3 = \{1, 2, 2, 4\}$ , which means that

$$\begin{aligned} \ell_0 &= 1, \\ \ell_1 &= 2, \\ \ell_2 &= 2, \\ \ell_3 &= 4. \end{aligned}$$

By formula (3.5),

$$\begin{aligned} b(3, 1) &= \ell_0 \ell_0 = 1 \cdot 1 = 1, \\ b(3, 2) &= \ell_0 \ell_1 + \ell_1 \ell_0 = 1 \cdot 2 + 2 \cdot 1 = 4, \\ b(3, 3) &= \ell_0 \ell_2 + \ell_1 \ell_1 + \ell_2 \ell_0 = 1 \cdot 2 + 2 \cdot 2 + 2 \cdot 1 = 8, \\ b(3, 4) &= \ell_0 \ell_3 + \ell_1 \ell_2 + \ell_2 \ell_1 + \ell_3 \ell_0 = 1 \cdot 4 + 2 \cdot 2 + 2 \cdot 2 + 4 \cdot 1 = 16, \\ b(3, 5) &= \ell_0 \ell_4 + \ell_1 \ell_3 + \ell_2 \ell_2 + \ell_3 \ell_1 + \ell_4 \ell_0 = 1 \cdot 0 + 2 \cdot 4 + 2 \cdot 2 + 4 \cdot 2 + 0 \cdot 1 = 20, \\ b(3, 6) &= \ell_0 \ell_5 + \ell_1 \ell_4 + \ell_2 \ell_3 + \ell_3 \ell_2 + \ell_4 \ell_1 + \ell_5 \ell_0 \\ &= 1 \cdot 0 + 2 \cdot 0 + 2 \cdot 4 + 4 \cdot 2 + 0 \cdot 2 + 0 \cdot 1 = 16, \\ b(3, 7) &= \ell_0 \ell_6 + \ell_1 \ell_5 + \ell_2 \ell_4 + \ell_3 \ell_3 + \ell_4 \ell_2 + \ell_5 \ell_1 + \ell_6 \ell_0 \\ &= 1 \cdot 0 + 2 \cdot 0 + 2 \cdot 0 + 4 \cdot 4 + 0 \cdot 2 + 0 \cdot 2 + 0 \cdot 1 = 16 \end{aligned}$$

which together with Eq. (3.6) yields

$$\begin{aligned} d(H_3, 1) &= 3[d(H_2, 1) + b(3, 1)] = 3[12 + 1] = 39, \\ d(H_3, 2) &= 3[d(H_2, 2) + b(3, 2)] = 3[12 + 4] = 48, \\ d(H_3, 3) &= 3[d(H_2, 3) + b(3, 3)] = 3[12 + 8] = 60, \\ d(H_3, 4) &= 3[d(H_2, 4) + b(3, 4)] = 3[0 + 16] = 48, \\ d(H_3, 5) &= 3[d(H_2, 5) + b(3, 5)] = 3[0 + 20] = 60, \\ d(H_3, 6) &= 3[d(H_2, 6) + b(3, 6)] = 3[0 + 16] = 48, \\ d(H_3, 7) &= 3[d(H_2, 7) + b(3, 7)] = 3[0 + 16] = 48 \end{aligned}$$

and finally,

$$H(H_3; \lambda) = 39\lambda + 48\lambda^2 + 60\lambda^3 + 48\lambda^4 + 60\lambda^5 + 48\lambda^6 + 48\lambda^7.$$

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