# A Refined Bound for the $Z_{1}$-Spectral Radius of Tensors 

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#### Abstract

A refined upper bound for the $Z_{1}$-spectral radius of tensors is given, which needs less computations than that presented by Wang et al. in [Applied Mathematics and Computation, 329 (2018) 266-277]. Numerical experiments involving Uniform distribution, Gaussian distribution, Poisson distribution and Binomial distribution are given to show the effectiveness of the proposed bound.


## 1. Introduction

The $Z_{1}$-eigenvalue of tensors and its corresponding eigenvectors are useful for computing the limiting probability distribution in high order Markov chain [1,10] and the PageRank vector in multilinear PageRank models [7,11], and also have applications in image matching [5], best rank-one approximation of tensors[14, 17], and hypergraph theory $[2,8]$.

Definition 1.1. [1] A real number $\lambda \in \mathbb{R}^{n}$ and a non-zero real vector $\mathbf{x}=\left(x_{1}, \ldots, x_{n}\right)^{\top} \in \mathbb{R}^{n}$ are called a $Z_{1^{-}}$ eigenvalue and a $Z_{1}$-eigenvector of an order m dimension $n$ real tensor $\mathcal{A}=\left(a_{i_{1}, \ldots, i_{m}}\right) \in \mathbb{R}^{[m, n]}\left(\mathbb{R}^{[m, n]}\right.$ denotes the set of the order $m$ dimension $n$ tensors over real numbers $\mathbb{R}$ ) if

$$
\begin{equation*}
\mathcal{A} \mathbf{x}^{m-1}=\lambda \mathbf{x},\|\mathbf{x}\|_{1}=\sum_{k=1}^{n}\left|x_{k}\right|=1 \tag{1}
\end{equation*}
$$

where $\mathcal{A} \mathbf{x}^{m-1}$ is a vector with its $i$-th component being

$$
\left(\mathcal{A x} x^{m-1}\right)_{i}=\sum_{i_{2}, \ldots, i_{m}=1}^{n} a_{i i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}, i \in[n]:=\{1, \ldots, n\}
$$

Furthermore, the $Z_{1}$-spectral radius of $\mathcal{A}$ is denoted by

$$
\rho_{z_{1}}(\mathcal{A})=\max \left\{|\lambda|: \lambda \in \sigma_{1}(\mathcal{A})\right\},
$$

where $\sigma_{1}(\mathcal{A})$ is the set of all $Z_{1}$-eigenvalues of $\mathcal{A}$.

[^0]There are a variety of results on the $Z_{1}$-eigenvalues and its corresponding $Z_{1}$-eigenvectors, such as, algorithms for computing $Z_{1}$-eigenvalues and its corresponding $Z_{1}$-eigenvectors [3], bounds for the $Z_{1}$ spectral radius $[9,12,16]$, and the uniqueness conditions for the positive $Z_{1}$-eigenvector for nonnegative tensors $[1,4,7,10,11]$.

Very recently, Wang et al. [16] provided an upper bound for the $Z_{1}$-spectral radius of tensors as follows.
Theorem 1.2. [16, Theorem 2.5] Let $\mathcal{A}=\left(a_{i_{1} i_{2} \cdots i_{m}}\right) \in \mathbb{R}^{[m, n]}$. Then

$$
\begin{equation*}
\rho_{z_{1}}(\mathcal{A}) \leq \min \left\{C_{1}(\mathcal{A}),(R(\mathcal{A}))^{\frac{1}{m-1}}\left(\min _{t \in[m] \backslash\{1\}} C_{t}(\mathcal{A})\right)^{\frac{m-2}{m-1}}\right\}, \tag{2}
\end{equation*}
$$

where $R(\mathcal{A}):=\max _{i \in[n]}\left\{r_{i}(\mathcal{A}):=\sum_{i_{2}, ., ., i_{m}=1}^{n}\left|a_{i i_{2}} \cdots i_{n}\right|\right\}$, and

$$
C_{t}(\mathcal{A}):=\max _{i_{s} \in[n], s \in[m] \backslash\{t\}} \sum_{i_{t}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{t} \cdots i_{m}}\right|, t \in[m] .
$$

As said in [16], if $m=2$, then the bound (2) reduces to the well-known Frobenius's bound [6] for the spectral radius $\rho(A)$ of a matrix $A$, i.e.,

$$
\rho(A) \leq \min \left\{C_{1}(A), C_{2}(A)\right\}
$$

where $C_{1}(A)$ and $C_{2}(A)$ are the maximum column sum and row sum of $\mathcal{A}$, respectively.
Although the bound (2) depends only on the entries of a given tensor $\mathcal{A}$, unlike matrices case it involves the term $R(\mathcal{A})$, and thus needs extra computations. In this paper, we give a refinement bound for the $Z_{1}$-spectral radius of tensors:

$$
\rho_{z_{1}}(\mathcal{A}) \leq \min _{t \in[m]} C_{t}(\mathcal{A})
$$

which has nothing to do with $R(\mathcal{A})$ like matrices case, and prove that the new bound is better than that in Theorem 1.2 ([16, Theorem 2.5]).

## 2. Main results

Let

$$
[n]^{m-1}=\left\{\left(i_{2}, i_{3}, \ldots, i_{m}\right): i_{j} \in[n], j=2,3, \ldots, m\right\} .
$$

Obviously, $[n]^{1}=[n]$.
Theorem 2.1. Let $\mathcal{A}=\left(a_{i_{1} i_{2} \cdots i_{m}}\right) \in \mathbb{R}^{[m, n]}$. Then

$$
\begin{equation*}
\rho_{z_{1}}(\mathcal{A}) \leq \min _{t \in[m]} C_{t}(\mathcal{A}) \tag{3}
\end{equation*}
$$

Proof. Suppose that a nonzero vector $\mathbf{x}=\left(x_{1}, x_{2}, \ldots, x_{n}\right)^{\top}$ with

$$
\|\mathbf{x}\|_{1}=\sum_{k=1}^{n}\left|x_{k}\right|=1
$$

such that $\mathcal{A} \mathbf{x}^{m-1}=\lambda \mathbf{x}$. We next consider the following two cases $t=1$ and $t=2, \ldots, m$.

Case I: $t=1$. From (1) we get

$$
\lambda x_{i_{1}}=\sum_{i_{2}, \ldots, i_{m}=1}^{n} a_{i_{1} i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}}, i_{1} \in[n] .
$$

Taking modulus in the above equation and using the triangle inequality give

$$
\begin{aligned}
|\lambda|=|\lambda| \sum_{i_{1}=1}^{n}\left|x_{i_{1}}\right| & \leq \sum_{i_{1}, i_{2}, \ldots, i_{m}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{m}}\right|\left|x_{i_{2}}\right| \cdots\left|x_{i_{m}}\right| \\
& =\sum_{i_{2}, \ldots, i_{m}=1}^{n}\left(\left|x_{i_{2}}\right| \cdots\left|x_{i_{m}}\right| \sum_{i_{1}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{m}}\right|\right) \\
& \leq\left(\sum_{i_{2}, \ldots, i_{m}=1}^{n}\left|x_{i_{2}}\right| \cdots\left|x_{i_{m}}\right| \max _{i_{s} \in[n],, \in[m] \backslash\{1\}} \sum_{i_{1}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{m}}\right|\right. \\
& =C_{1}(\mathcal{A}),
\end{aligned}
$$

where the last equality holds because

$$
\sum_{i_{2}, \ldots, i_{m}=1}^{n}\left|x_{i_{2}}\right| \cdots\left|x_{i_{m}}\right|=\prod_{k=2,3, \ldots, m}\left(\sum_{i_{k}=1}^{n}\left|x_{i_{k}}\right|\right)=1
$$

Thus, $\rho_{z_{1}}(\mathcal{A}) \leq C_{1}(\mathcal{F})$.
Case II: $t=2, \ldots, m$. Let $\left|x_{k}\right|=\max _{i \in[n]}\left|x_{i}\right|$. Then $\left|x_{k}\right| \neq 0$. From the $k$-th equality of (1) we get

$$
\lambda x_{k}=\sum_{\left(i_{2}, \ldots, i_{m}\right) \in[n]^{m-1}} a_{k i_{2} \cdots i_{m}} x_{i_{2}} \cdots x_{i_{m}} .
$$

Taking modulus in the above equation and using the triangle inequality give

$$
\begin{aligned}
\left|\lambda \| x_{k}\right| & \leq \sum_{\left(i_{2}, \ldots, i_{m}\right) \in[n]^{m-1}}\left|a_{k i_{2} \cdots \cdots i_{m}}\right|\left|x_{i_{2}}\right| \cdots\left|x_{i_{m}}\right| \\
& \left.=\sum_{i_{p}=1}^{n}\left(\sum_{\left(i_{2}^{\prime}, \ldots, i_{m-1}^{\prime}\right) \in[n]^{m-2}}\left|a_{k i_{2}^{\prime} \cdots i_{p} \cdots i_{m-1}^{\prime}}\right| \prod_{\substack{s=2,2 \\
s \neq p}}^{m-1}\left|x_{i_{s}}^{\prime}\right|\right)\left|x_{i_{p}}\right|\right) \\
& \leq\left(\max _{i_{p} \in[n]}\left(\sum_{\left(i_{2}^{\prime}, \ldots, i_{m-1}^{\prime}\right) \in[n]^{m-2}}\left|a_{k i_{2}^{\prime} \cdots \cdots i_{p} \cdots i_{m-1}^{\prime}}\right| \prod_{\substack{s=2, p \\
s \neq p}}^{m-1}\left|x_{i_{s}}^{\prime}\right|\right) \sum_{i_{p}=1}^{n}\left|x_{i_{p}}\right|\right. \\
& =\max _{i_{p} \in[n]}\left(\sum_{\left(i_{2}^{\prime}, \ldots, i_{m-1}^{\prime}\right) \in[n]^{m-2}}\left|a_{k i_{2}^{\prime} \cdots i_{p} \ldots i_{m-1}^{\prime}}\right| \prod_{\substack{s=2,2 \\
s \neq p}}^{m-1}\left|x_{i_{s}^{\prime}}^{\prime}\right|\right)
\end{aligned}
$$

$$
\begin{aligned}
& =\max _{i_{p} \in[n]}\left(\sum_{i_{q}=1}^{n}\left(\sum_{\left(i_{2}^{\prime \prime}, \ldots, i_{m-2}^{\prime \prime}\right) \in[n]^{m-3}}\left|a_{k i_{2}^{\prime \prime} \ldots i_{p} \cdots i_{q} \cdots i_{m-2}^{\prime \prime}}\right| \prod_{\substack{s=2, s f p, q}}^{m-2}\left|x_{i_{s}}^{\prime \prime}\right|\right)\left|x_{i_{q}}\right|\right) \\
& \leq \max _{i_{p} \in[n]}\left(\left(\max _{i_{q} \in[n]}\left(\sum_{\left(i_{2}^{\prime \prime}, \ldots, i_{m-2}^{\prime \prime}\right) \in[n]^{m-3}}\left|a_{k i_{2}^{\prime \prime} \ldots i_{p} \cdots i_{q} \cdots i_{m-2}^{\prime \prime}}\right| \prod_{\substack{s=2, s \neq p, q}}^{m-2}\left|x_{i_{s}}^{\prime \prime \prime}\right|\right) \sum_{i_{q}=1}^{n}\left|x_{i_{q}}\right|\right)\right. \\
& =\max _{i_{p} \in[n]}\left(\max _{i_{q} \in[n]}\left(\sum_{\left(i_{i}^{\prime \prime}, \ldots, i_{m-2}^{\prime \prime}\right) \in[n]^{n-3}}\left|a_{k i_{2}^{\prime \prime} \cdots i_{p} \cdots i_{q} \cdots \cdots i_{m-2}^{\prime \prime}}\right| \prod_{\substack{s=2, q \\
s p p, q}}^{m-2}\left|x_{i_{s}^{\prime \prime}}\right|\right)\right) \\
& =\max _{\left(i_{p}, i_{q}\right) \in[n]^{2}}\left(\sum_{\left(i_{2}^{\prime \prime}, \ldots, i_{m-2}^{\prime \prime}\right) \in[n]^{m-3}}\left|a_{k i_{2}^{\prime \prime} \cdots i_{p} \cdots i_{q} \cdots i_{m-2}^{\prime \prime}}\right| \prod_{\substack{s=2, t \\
s \neq p, q}}^{m-2}\left|x_{i_{s}^{\prime \prime}}\right|\right) \\
& =\max _{\left(i_{2}^{*}, \ldots, i_{m-1}^{i_{m-1}}\right) \in[n]^{m-2}}\left(\sum_{i_{t}=1}^{n}\left|a_{k_{2}^{*} \ldots i_{i} \cdots i_{m-1}^{*}} \| x_{i_{t}}\right|\right) \\
& \leq\left(\max _{\left(i_{2}^{*}, \ldots, i_{m-1}^{i_{2}^{*}}\right) \in[n]^{m-2}}\left(\sum_{i_{t}=1}^{n}\left|a_{k i_{2}^{*} \cdots \cdots i i_{1} \cdots i_{m-1}^{*}}\right|\right)\right)\left|x_{k}\right|
\end{aligned}
$$

Dividing $\left|x_{k}\right| \neq 0$ on both sides yields

$$
\begin{aligned}
|\lambda| & \leq \max _{\left(i_{2}^{*}, \ldots, i_{m-1}^{*}\right) \in[n]^{m-2}}\left(\sum_{i_{t}=1}^{n}\left|a_{k i_{2}^{*} \cdots i_{t} \cdots i_{m-1}^{*}}\right|\right) \\
& \leq \max _{i_{1} \in[n]} \max _{\left(i_{2}^{*}, \ldots, i_{m-1}^{*}\right) \in[n]^{m-2}}\left(\sum_{i_{t}=1}^{n}\left|a_{i_{1} i_{2}^{*} \cdots i_{t} \cdots i_{m-1}^{*}}\right|\right) \\
& =\max _{i_{s} \in[n], s \in[m] \backslash\{t\rangle}\left(\sum_{i_{t}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{t} \cdots i_{m}}\right|\right) .
\end{aligned}
$$

Apparently, the inequality above holds for any $t=2, \ldots, m$, and hence

$$
|\lambda| \leq \min _{t \in[m] \backslash\{1\}} \max _{i_{s} \in[n], s \in[m] \backslash\{t\rangle}\left(\sum_{i_{t}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{t} \cdots i_{m}}\right|\right)=\min _{t \in[m] \backslash\{1\}} C_{t}(\mathcal{A}),
$$

consequently,

$$
\rho_{z_{1}}(\mathcal{A}) \leq \min _{t \in[m] \backslash\{1\}} C_{t}(\mathcal{A}) .
$$

The conclusion follows from Case I and Case II.
If $\mathcal{A} \in \mathbb{R}^{[m, n]}$ is a nonnegative tensor, then the bound (3) reduces to

$$
\rho_{z_{1}}(\mathcal{F}) \leq \min _{t \in[m]} \max _{i_{s} \in[n], s \in[m] \backslash\{t\}} \sum_{i_{t}=1}^{n} a_{i_{1} i_{2} \cdots i_{H} \cdots i_{m}},
$$

which is the exact upper bound in Corollary 3.6 of [9] for the weakly symmetric nonnegative irreducible tensor case. Apparently, the bound (3) needs less computations than the bound (2) because the latter has to compute $R(\mathcal{A})$. Next, we establish a comparison result to show that the bound (3) is less than or equal to the bound (2).

Theorem 2.2. Let $\mathcal{A}=\left(a_{i_{1} i_{2} \cdots i_{m}}\right) \in \mathbb{R}^{[m, n]}$. Then

$$
\min _{t \in[m]} C_{t}(\mathcal{A}) \leq \min \left\{C_{1}(\mathcal{A}),(R(\mathcal{A}))^{\frac{1}{m-1}}\left(\min _{t \in[m] \backslash 1\}} C_{t}(\mathcal{A})\right)^{\frac{m-2}{m-1}}\right\},
$$

where $R(\mathcal{A})$ and $C_{t}(\mathcal{A}), t \in[m]$ are defined as in Theorem 1.2.
Proof. Note that for any $t=2,3, \ldots, m$,

$$
\max _{i_{s} \in[n], s \in[m] \backslash\{t\}} \sum_{i_{t}=1}^{n}\left|a_{i_{1} i_{2} \cdots i_{t} \cdots i_{m}}\right| \leq \max _{i \in[n]} r_{i}(\mathcal{A}) .
$$

Hence, $\min _{t \in[m]} C_{t}(\mathcal{A}) \leq \min _{t \in[m] \backslash 11} C_{t}(\mathcal{A}) \leq R(\mathcal{A})$. Furthermore, from $\min _{t \in[m]} C_{t}(\mathcal{A}) \leq C_{1}(\mathcal{A})$, we have

$$
\begin{aligned}
\min _{t \in[m]} C_{t}(\mathcal{A}) & =\min \left\{C_{1}(\mathcal{A}), \min _{t \in[m]} C_{t}(\mathcal{A})\right\} \\
& \leq \min \left\{C_{1}(\mathcal{A}),\left(\min _{t \in[m]} C_{t}(\mathcal{A})\right)^{\frac{1}{m-1}}\left(\min _{t \in[m] \backslash\{1\}} C_{t}(\mathcal{A})\right)^{\frac{m-2}{m-1}}\right\} \\
& \leq \min \left\{C_{1}(\mathcal{A}),(R(\mathcal{A}))^{\frac{1}{m-1}}\left(\min _{t \in[m] \backslash\{1\}} C_{t}(\mathcal{A})\right)^{\frac{m-2}{m-1}}\right\}
\end{aligned}
$$

The proof is complete.
Remark here that besides the bound (2) in Theorem 1.2 ([16, Theorem 2.5]), there are another bounds for the $Z_{1}$-spectral radius, for instance, in 2015, Li et al. [9, Theorem 2.1] derived the following upper bound about the $Z_{1}$-spectral radius of $\mathcal{A}$ :

$$
\rho_{z_{1}}(\mathcal{A}) \leq \min _{k \in[m]} \max _{i_{k} \in[n]} \sum_{\substack{\left.\left.i_{s} \in[n]\right] \\ s \in[m] \mid k k\right]}}\left|a_{i_{1} \cdots i_{k} \cdots i_{m}}\right| .
$$

As stated in [16, Remark 3],

$$
\min \left\{C_{1}(\mathcal{A}),(R(\mathcal{A}))^{\frac{1}{m-1}}\left(\min _{t \in[m] \backslash 1\}} C_{t}(\mathcal{A})\right)^{\frac{m-2}{m-1}}\right\} \leq \min _{k \in[m]} \max _{i_{k} \in[n]} \sum_{\substack{\left.i_{k} \in[n]\right] \\ s \in[m] \mid(k]}}\left|a_{i_{1} \cdots i_{k} \cdots i_{m}}\right| .
$$

Hence,

$$
\min _{t \in[m]} C_{t}(\mathcal{A}) \leq \min _{k \in[m]} \max _{i_{k} \in[n]} \sum_{\substack{\left.\left.i_{\in} \in[n]\right] \\ s \in[m]|k|\right]}} \mid a_{i_{1} \cdots i_{k} \cdots i_{m} \mid .} .
$$

This implies that the bound in Theorem 2.1 is better than that in [9, Theorem 2.1].


Figure 1: The bound differences for four distributions entries.

Example 2.3. Consider $4 \times 10^{3}$ order 4 dimensional 2 tensors generated by the way from [16], i.e., tensors are implemented randomly with four different distributions (Uniform distribution, Gaussian distribution, Poisson distribution and binomial distribution) entries. In uniform distribution case, all entries are in the range of $[0,1]$. In gaussian distribution case, the parameters $\mu$ and $\sigma$ are generated randomly in the range of $[0,1]$. For convenience, all the entries of tensor $\mathcal{A}$ are shifted to be positive. In poisson distribution case, the parameter $\lambda$ is set to be 10 . In binomial distribution case, the number of entries is set to be 100. And the probability of success for each trial p is set to be 0.5 .

The differences of the bounds in Theorem 1.2, Theorem 2.1 and [9, Theorem 2.1] are drawn in Figure 1, where the star symbol in red color ' $*$ ' means the upper bound in [9, Theorem 2.1] minus the upper bound in Theorem 2.1, and the cross symbol in blue color + means the upper bound in Theorem 1.2 minus the upper bound in Theorem 2.1. From all sub-figures it is easy to see that there are no ' $*$ ' and ' + ' below zero. This means that the upper bound in Theorem 2.1 is better than that in Theorem 1.2 and [9, Theorem 2.1].

## 3. Conclusions

In this paper, we give a new upper bound for the $Z_{1}$-spectral radius for tensors, and it needs less computations, and is sharper than that in [16].

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